

Classes of Errors
BIOS Error-Handling Functions
Practical Error Handling
Character I/O Errors
Disk Errors
Improving Error Messages



Dealing with Hardware Errors

This chapter describes the enhancements you can make to improve CP/M's somewhat primitive error handling. It covers the general classes of errors that the BIOS may have to handle. It describes some of the underlying philosophical aspects of errors, how to detect them, and how to correct them or otherwise make the best of the situation.

At the end of the chapter are some example error-handling subroutines. Some of these have already been shown in the previous chapter as part of the enhanced BIOS (Figure 8-10); they are repeated here so that you can see them in isolation.

Classes of Errors

Basically, the user perceives only two classes of errors — those that are user-correctable and those that are not. There is a third, almost invisible class of errors—those that are recoverable by the hardware or software without the user's intervention.

The possible sources for hardware errors vary wildly from one computer system to another, since error detection is heavily dependent on the particular logic in the hardware. The BIOS can detect some hardware-related errors — mainly errors caused when something takes too long to happen, such as when a recalcitrant printer does not react in a specified length of time.

The BDOS has no built-in hardware detection code. It can detect *system* errors, such as an attempt to write to a disk file that is marked “Read-Only” in the file directory or attempts to access files that are not on the disk. These BDOS-detected errors, however, generally are unrelated to the well-being of the hardware. For example, a disk controller with a hardware problem could easily overwrite a sector of the directory, thereby deleting several files. This error would not show up until the user tried to use one of the now-departed files.

BIOS Error-Handling Functions

The error-handling code in the BIOS has to serve the following functions:

- Detection
- Analysis
- Indication
- Correction.

Error Detection

Clearly, before any later steps can be taken, an error must be detected. This can be done by the software alone or by the BIOS interacting with error-detecting logic in the hardware. In general, the only errors that the BIOS can detect unassisted are caused when certain operations take longer to complete than expected. Because the writer of the BIOS knows the operating environment of the specific peripherals in the system, the code can predict how long a particular operation should take and can signal an error when this time is exceeded. This would include such problems as printers that fail to react within a specified time period.

The BIOS can work in cooperation with the hardware to determine whether the hardware itself has detected an error. Armed with the hardware's specifications, the BIOS can input information on controller or device status to trigger error-detecting logic. How this should be done depends heavily on the peripheral devices in your computer system and the degree to which these devices have “smart” controllers capable of processing independently of the computer. Unfortunately, many manufacturers document the significance of individual status bits that indicate errors, but not combinations of errors, or what to do when a particular error occurs.

Error Analysis

Given that your BIOS has detected an error, it must first determine the class of error; that is, whether or not the error can be corrected by simply trying the operation again. Some errors appear at first to be correctable, but retrying the operation several times still fails to complete it. An example would be a check-sum error while reading a disk sector. If several attempts to read the sector all yield an error, then it becomes a “fatal” error. The code in your BIOS must be capable of initial classification and then subsequent reclassification if remedial action fails.

Other types of errors can be classified immediately as fatal errors—nothing can be done to save the situation. For example, if the floppy disk controller indicates that it cannot find a particular sector number on a diskette (due to an error in formatting), there is nothing that the BIOS can do other than inform the user of the problem and supply other helpful information.

Analysis of errors may require some basic research, such as inducing failures in the hardware and observing combinations of error indicators. For example, some printers (interfaced via a parallel port) indicate that they are “Out of Paper” or “Busy” when, in fact, they are switched off. The BIOS should detect this condition and tell the user to switch the printer on, not load more paper.

Error Indication

An incomplete or cryptic error message is infuriating. It is the functional equivalent of saying, “There has been an error. See if you can guess what went wrong!”

An error message, to be complete, should inform the recipient of the following:

- The fact that an error has occurred.
- Whether or not automatic recovery has been attempted and failed.
- The details of the error, if need be in technical terms to assist a hardware engineer.
- What possible choices the user has now.

To put these points into focus, consider the error message that can be output by CP/M after you have attempted to load a program by entering its name into the CCP. What you see on the console is the following dialog:

```
A>myprog<cr>
BAD LOAD
A>
```

All you know is that there has been an error, and you must guess what it is, even though the specific cause of the error was known to CP/M when it output the message. This error message is output by the CCP when it attempts to load a

“.COM” file larger than the current transient program area. The message “BAD LOAD” is only understandable *after* you know what the error is. Even then, it does not tell you what went wrong, whether there is anything you can do about it, and how to go about doing it.

To be complete, this error message could say something like this:

```
A>myprog<cr>
"MYPROG.COM" exceeds the available memory space by
1,024 bytes, and therefore cannot be loaded under the
current version of CP/M.
```

Notice how the message tells you what the problem is, and even quantifies it so that you can determine its severity (you need to get 1K more memory or reduce the program's size). It also tells you how you stand—you cannot load this program under the current version of CP/M, so retrying the operation is futile.

Not many systems programmers like to output messages like the example above. They argue that such a message is too long and too much work for something that does not happen often. Admittedly, the message *is* too long. It could be shortened to read

```
(131) Program 1,024 bytes too large to load.
```

This conveys the same information; the number in parentheses can serve as a reference to a manual where the full impact of the message should be described.

The major problem with the way error messages are designed is that they usually are written by programmers to be read by nontechnical lay users, and programmers are notoriously bad at guessing what nonexperts need to know.

Error indications you design should address the following issues, from the point of view of the user:

- The cause of the error
- The severity of the error
- The corrective action that has and can be taken.

Examine the error messages in the error processor for the example BIOS in Figure 8-10, from line 03600 onward. Although these are an improvement on the BDOS all-purpose

```
BDOS Error on A: Bad Sector
```

even these messages do not really meet all of the requirements of a good error message system.

Another often overlooked aspect of errors is that most hardware errors form a pattern. This pattern is normally only discernible to the trained eye of a hardware maintenance engineer. When these engineers are called to investigate a problem,

they will quiz the user to determine whether a given failure is an isolated incident or part of an ongoing pattern. This is why an error message should contain additional technical details. For example, a disk error message should include the track and sector used in the operation that resulted in an error. Only with these details can the engineer piece together the context of a failure or group of failures.

Error Correction

Given that a lucid error message has been displayed on the console, the user is still confronted with the question: “Now what do I do?” Not only can this be difficult for the user to answer, but also the particular solution decided upon can be hard for the BIOS to execute.

Normally, there are three possible options in response to errors:

- Try the operation again
- Ignore the error and attempt to continue
- Abort the program causing the error and return to CP/M.

For some errors, retrying can be effective. For example, if you forget to put the printer on-line and get a “Printer Timeout” error message, it is easy to put the printer back on-line and ask the BIOS to try again to send data to the printer.

Seldom can you ignore an error and hope to get sensible results from the machine; many disk controllers do not even transfer data between themselves and the disk drive if an error has been detected. Only ignorant users, or brave ones in desperation, ignore errors.

Aborting the program causing the error is a drastic measure, although it does escape from what could otherwise be a “deadly embrace” situation. For example, if you misassign the printer to an inactive serial port and turn on printer echoing (with the CONTROL-P toggle), you will send the system into an endless series of “Printer Timeout” messages. If you abort the program, the error handler in the BIOS executes a System Reset function (function 0) in the BDOS, CP/M warm boots, and control is returned to the CCP. In the process, the printer toggle is reset and the circle is broken.

Practical Error Handling

This section discusses several errors, describing their causes and the way in which the BIOS and the user can handle them when they occur.

Character I/O Errors

At the BIOS level, most detectable errors related to character input or output will be found by the hardware chips.

Parity Error

Parity, in this context, refers to the number of bits set to 1 in an 8-bit character. The otherwise unused eighth bit in ASCII characters can be set to make this number always odd, or alternatively, always even. Your computer hardware can be programmed to count the number of 1 bits in each character and to generate an error if the number is odd (odd parity) or, alternatively, if it is even (even parity). If the hardware on the other end of the line is programmed to operate in the same mode, parity checking provides a primitive error-detection mechanism — you can tell that a character is bad, but not what it should have been.

CP/M does not provide a standard mechanism for reporting a parity error, so your only option is to reset the hardware and substitute an ASCII DEL (7FH; delete) character in the place of the erroneous character.

If your BIOS is operating in a highly specialized environment, you may need to count the number of such parity errors so that a utility program can report on the overall performance of the system.

Framing Error

When an 8-bit ASCII character is transmitted over a serial line, the eight bits are transmitted serially, one after the other. A *start* bit is transmitted first, followed by the data character and then a *stop* bit. If the hardware fails to find the stop and start bits in the correct positions, a *framing error* will occur. Again, the only option available to the BIOS is to reset the hardware chip and substitute an ASCII DEL.

Overrun Error

This error occurs when incoming data characters arrive faster than the program can handle them, so that the last characters overrun those being processed by the hardware chip. This error can normally be avoided by the use of serial line protocols, such as those in the example BIOS in Figure 8-10.

An *overrun error* implies that the protocol has broken down. As with the parity and framing errors, almost the only option is to reset the hardware and substitute a DEL character.

Printer Timeout Error

This is one of the few errors where the BIOS can sensibly attempt an error recovery. The error occurs when the BIOS tries to output a character to a serial printer and finds that the printer is not ready for more than, say, 30 seconds. The most common cause of this error is that the user forgets to put the printer on-line. Many printers require that they be off-line during a manual form feed, and users will often forget to push the on-line button afterward.

After a 30-second delay, the BIOS can send a message to the console device(s) informing the user of the error and asking the user to choose the appropriate course of action. Note that console output can be directed to more than one device.

Parallel Printers

Printers connected to your system by means of a parallel port can indicate their status to the computer much more easily than can serial printers. They can communicate such error states as “Out of Paper,” “End of Ribbon,” and “Off-line.”

These single-error indicators can also be used in combination to indicate whether the printer cable is connected, or even whether the printer is receiving power. You need to experiment, deliberately putting the printer into these states and reading status in order to identify them. It is misleading to indicate to the inexperienced user that the printer is “Out of Paper” when the problem is that the data cable has inadvertently become disconnected.

However, each of these errors can be dealt with in the same way as the serial printer’s timeout problem: display an error message and request the user’s choice of action.

Example Printer Error Routine

Figure 9-1 shows an example of a program that handles printer errors. It consists of several subroutines, including

- The error detection classification and indication routine
- The error correction routine.

It uses other subroutines that are omitted from the figure to avoid obscuring the logic. These subroutines are listed in full in the example BIOS in Figure 8-10.

```

;      This example shows, in outline form, how to handle the
;      situation when a serial printer remains busy for too long.
;      It is intended that this generic example show how to
;      deal with this class of errors.
;
;      The example presupposes the existence of a clock interrupt
;      every 16.666 milliseconds (1/60th of a second), and that
;      control will be transferred to the Real Time Clock service
;      routine each time the clock "ticks".
;
;      Figure 8-10 shows a more complete example, installed in a real
;      BIOS.
;
0000 =  B$System$Reset      EQU    0      ;BDOS system reset function
0005 =  BDOS                EQU    5      ;BDOS entry point
;
0000 00  Printer$Timeout$Flag:  DB    0      ;This flag is set by the interrupt
;      service subroutine that is called
;      when the watchdog timer subroutine
;      count hits zero (after having
;      counted down a 30-second delay)

0708 =  Printer$Delay$Count  EQU    1800  ;Given a clock period of 16.666 ms
;      this represents a delay of 30 secs

```

Figure 9-1. Serial printer error handling

```

000D =      ;
000A =      CR          EQU    ODH    ;Carriage return
          LF          EQU    OAH    ;Line feed
          ;
          Printer$Busy$Message:
0001 0D0A      DB      CR,LF
0003 507269E74 DB      'Printer has been busy for too long.',CR,LF
0028 436865636B DB      'Check that it is on-line and ready.',CR,LF,0
          ;
004E 00      Printer$Character:  DB      0      ;Save area for the data character
          ;                                     ; to be output
          ;
LIST:      ;<=== Main BIOS entry point
          ;<=== I/O redirection code occurs here
          ;.....

004F 79      MOV      A,C          ;Save the data character
0050 324E00   STA      Printer$Character

Printer$Retry:
0053 010807   LXI      B,Printer$Delay$Count ;This is the count of the number
          ; of clock ticks before the watchdog
          ; subroutine call
0056 217E00   LXI      H,Printer$Timed$Out    ; <= this address
0059 CDA300   CALL     Set$Watchdog          ;Sets the watchdog running

Printer$Wait:
005C CDA300   CALL     Get$Printer$Status    ;See if the printer is ready to
          ; accept a character for output
          ; This includes checking if the printer
          ; is "Busy" because the driver is
          ; waiting for XON, ACK, or DTR to
          ; come high
005F C26C00   JNZ     Printer$Ready         ;The printer is now ready

0062 3A0000   LDA     Printer$Timeout$Flag  ;Check if the watchdog timer has
          ; hit zero (if it does, the
          ; watchdog routine will call
          ; the Printer$Timed$Out code
          ; that sets this flag)

0065 B7      ORA     A
0066 C28400   JNZ     Display$Busy$Message  ;Yes, so display message to
          ; indicate an error has occurred
0069 C35C00   JMP     Printer$Wait          ;Otherwise, check if printer is
          ; now not busy

Printer$Ready:
          ;The printer is now ready to output
          ; a character, but before doing so,
          ; the watchdog timer must be reset
          ;Ensure no false timeout occurs
006C F3      DI
006D 010000   LXI     B,0
0070 CDA300   CALL   Set$Watchdog          ;This is done by setting the count
          ; to zero
0073 FB      EI

0074 3A4E00   LDA     Printer$Character      ;Get character to output
0077 11A300   LXI     D,Printer$Device$Table ;DE -> device table for printer
007A CDA300   CALL   Output$Data$Byte       ;Output the character to the printer

007D C9      RET                    ;Return to the BIOS's caller
          ;
          ;
Printer$Timed$Out:
          ;Control arrives here from the
          ; watchdog routine if the
          ; watchdog count ever hits zero
          ; This is an interrupt service
          ; routine
          ;All registers have been saved
          ; before control arrives here
          ;Set printer timeout flag
007E 3EFF      MVI     A,OFFH
0080 320000   STA     Printer$Timeout$Flag  ;Set printer timeout flag
0083 C9      RET                    ;Return back to the watchdog

          ;Interrupt service routine

```

Figure 9-1. (Continued)


```

;
; Display$Busy$Message:
;Printer has been busy for
; 30 seconds or more
;Reset timeout flag
0084 AF          XRA    A
0085 320000     STA    Printer$Timeout$Flag

0088 210100     LXI    H,Printer$Busy$Message ;Output error message
008B CDA300     CALL   Output$error$Message

008E CDA300     CALL   Request$user$Choice ;Displays a Retry, Abort, Ignore?
; prompt, accepts a character from
; the keyboard, and returns with the
; character, converted to upper
; case in the A register
;Check if Retry

0091 FE52       CPI    'R'
0093 CA5300     JZ     Printer$Retry
0096 FE41       CPI    'A' ;Check if Abort
0098 CA9E00     JZ     Printer$Abort
009B FE49       CPI    'I' ;Check if Ignore
009D C8         RZ

;
; Printer$Abort:
009E 0E00       MVI    C,B$System$Reset ;Issue system reset
00A0 C30500     JMP    BDOS ;No need to give call as
; control will not be returned

;
;
; Dummy subroutines
; These are shown in full in Figure 8-10. The line numbers in
; Figure 8-10 are shown in the comment field below
;
Printer$Device$Table: ;Line 01300 (example layout)
Request$user$Choice: ;Line 03400
Output$error$Message: ;Line 03500
Get$Printer$Status: ;Line 03900 (similar code)
Output$data$Byte: ;Line 05400 (similar code)
Set$Watchdog: ;Line 05800

```

Figure 9-1. Serial printer error handling (continued)

Disk Errors

Disks are much more complicated than character I/O devices. Errors are possible in the electronics and in the disk medium itself. Most of the errors concerned with electronics need only be reported in enough detail to give a maintenance engineer information about the problem. This kind of error is rarely correctable by retrying the operation. In contrast, media errors often can be remedied by retrying the operation or by special error processing software built into the BIOS. This chapter discusses this class of errors.

Media errors occur when the BIOS tries to read a sector from the disk and the hardware detects a check-sum failure in the data. This is known as a *cyclical redundancy check* (CRC) error. Some disk controllers execute a read-after-write check, so a CRC error can also occur during an attempt to write a sector to the disk.

With floppy diskettes, the disk driver should retry the operation at least ten times before reporting the error to the user. Then, because diskettes are inexpensive and replaceable, the user can choose to discard the diskette and continue with a new one.

With hard disks, the media cannot be exchanged. The only way of dealing with bad sectors is to replace them logically, substituting other sectors in their place.

There are two fundamentally different ways of doing this. Figure 9-2 shows the scheme known as sector sparing—substituting sectors on an outer track for a sector that is bad.

The advantage of this scheme is that it is dynamic. If a sector is found to be bad in a read-after-write check, even after several retries, then the data intended for the failing sector can be written to a spare sector. The failing sector's number is placed into a spare-sector directory on the disk. Thereafter, the disk drivers will be redirected to the spare sector every time an attempt is made to read or write the bad sector.

The disadvantage of this system is that the read/write heads on the disk must move out to the spare sector and then back to access the next sector. This can be a problem if you attempt to make a high-speed backup on a streaming tape drive (one that writes data to a tape in a single stream rather than in discrete blocks). The delay caused by reading the spare sector interrupts the data flow to the streaming tape drive.

You need a special utility program to manipulate the spare-sector directory, both to substitute for a failing sector manually and to attempt to rewrite a spare sector back onto the bad sector.

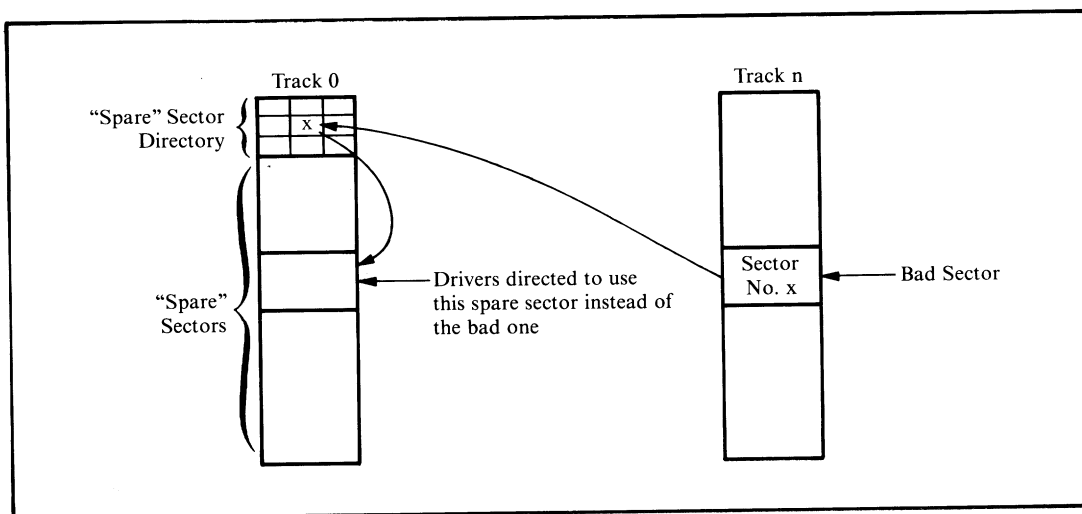


Figure 9-2. Sector sparing

Figure 9-3 shows another scheme for dealing with bad sectors. In this method, bad sectors are skipped rather than having sectors substituted for them.

The advantage of sector skipping is that the heads do not have to perform any long seeks. The failing sector is skipped, and the next sector is used in its place. Because of this, sector skipping can give much better performance. Data can be read off the disk fast enough to keep a streaming tape drive “fed” with data.

The disadvantage of sector skipping is that it does not lend itself to dynamic operation. The bad sector table is best built during formatting. Once data has been written to the disk, if a sector goes bad, all subsequent sectors on the disk must be “moved down one” to make space to skip the bad sector. On a large hard disk, this could take several minutes.

Example Bad Sector Management

Sector sparing and sector skipping use similar logic. Both require a spare-sector directory on each physical disk, containing the sector numbers of the bad sectors. This directory is read into memory during cold start initialization. Thereafter, all disk read and write operations refer to the memory-resident table to see if they are about to access a bad sector.

For sector sparing, if the sector about to be read or written is found in the spare directory, its position in the directory determines which spare sector should be read.

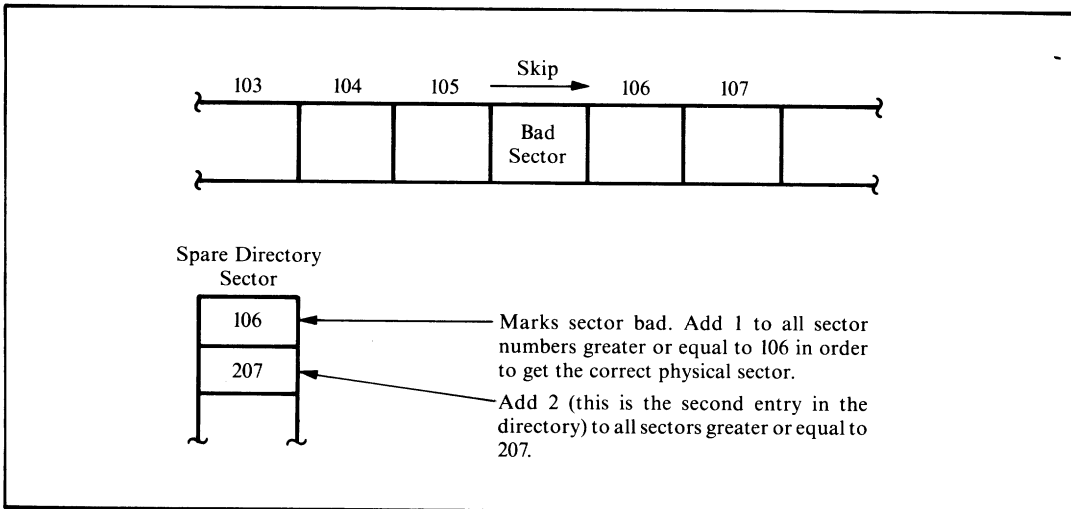


Figure 9-3. Sector skipping

In the case of sector skipping, every access to the disk makes the driver check the bad sector directory. The directory is used to tell how many bad sectors exist between the start of the disk and the failing bad sector. This number must be added to the requested track and sector to compensate for all the bad sectors.

The physical low-level drivers need four entry points:

- Read the specified sector without using bad sector management. This is used to read in the spare directory itself.
- Write the specified sector without using bad sector management. This is used to write the spare directory onto the disk, both to initialize it and to update it.
- Read and write the sector using bad sector management. These entry points are used for normal disk input/output.

Figure 9-4 shows the code necessary for both sector sparing and (using conditional code) sector skipping.

```

;      This example shows the modifications to be made in order
;      to implement bad sector management using sector sparing
;      and sector skipping.
;
0000 = False      EQU      0
FFFF = True       EQU      Not False
;
0000 = Sector$Sparing EQU    False
FFFF = Sector$Skipping EQU    Not Sector$Sparing
;
;      Additional equates and definitions
;
Spare$Directories:      ;Table of spare directory addresses
;Note: The directories themselves
; are declared at the end of the
; BIOS
0000 D500              DW      Spare$Directory$0      ;Physical disk 0
0002 9701              DW      Spare$Directory$1      ;Physical disk 1
;
Spare$Dir$In$Memory:    ;Flags used to indicate whether spare
0004 00                DB      0                    ; directory for a given physical disk
0005 00                DB      0                    ; has been loaded into memory. Set by SELDSK
;
;
0000 = Spare$Track      EQU      0                    ;Track containing spare directory
; sectors
0004 = Spare$Sector     EQU      4                    ;Sector containing directory
0005 = First$Spare$Sector EQU    Spare$Sector + 1
;
;      Variables set by SELDSK
;
Selected$Spare$Directory:
0006 0000              DW      0                    ;Pointer to directory
0008 00                DB      0                    ;Logical disk number
0009 00                DB      0                    ;Floppy/hard disks
000A 00                DB      0                    ;Deblocking flag
000B 00                DB      0                    ;Physical disk number
;
000C 0000              DW      0                    ;) These variables are part of the command
000E 00                DB      0                    ;) block handed over to the disk controller

```

Figure 9-4. Bad sector management

```

;
8000 = Maximum$Track EQU 32768 ;Used as a terminator
0012 = Sectors$Per$Track EQU 18
0000 = First$Sector$On$Track EQU 0
;
;
Disk$Parameter$Headers:
;-----
;Standard DPH Declarations
;-----
;
;
; Equates for disk parameter block
;
; The special disk parameter byte that precedes each disk
; parameter block, needs to be rearranged so that a
; physical disk drive number can be added.
;
;
; Disk types
;
; vvvv--- Physical disk number
0010 = Floppy$5 EQU 0$001$0000B ;5 1/4" mini floppy
0020 = Floppy$8 EQU 0$010$0000B ;8" floppy (SS SD)
0030 = M$Disk EQU 0$011$0000B ;Memory disk
0040 = H$Disk$10 EQU 0$100$0000B ;Hard disk - 10 megabyte
;
0070 = Disk$Type$$Mask EQU 0$111$0000B ;Masks to isolate values
000F = Physical$Disk$Mask EQU 0$000$1111B
;
; Blocking/deblocking indicator
;
0080 = Need$Deblocking EQU 1$000$0000B ;Sector size > 128 bytes
;
;
; Disk parameter blocks
;
;-----
; Standard DPB's for A: and B:
;-----
;
;-----
; Logical disk C:
; Extra byte indicates disk type
; deblocking requirements and physical
; disk drive.
000F C0 DB H$Disk$10 + Need$Deblocking + 0 ; Physical drive 0
Hard$5$Parameter$Block$C:
;-----
; Standard format parameter block
;-----
;
;
0010 C0 DB H$Disk$10 + Need$Deblocking + 0 ; Physical drive 0
Hard$5$Parameter$Block$D:
;-----
; Standard format parameter block
;-----
;
;
0004 = NumberOf$Logical$Disks EQU 4
;
;
SELDSK: ;Select disk in register C
; C = 0 for drive A, 1 for B, etc.
; Return the address of the appropriate
; disk parameter header in HL, or 0000H
; if the selected disk does not exist.
;
0011 21000 LXI H,0 ;Assume an error
0014 79 MOV A,C ;Check if requested disk valid
;
0015 FE04 CPI NumberOf$Logical$Disks
0017 D0 RNC ;Return if > maximum number of disks

```

Figure 9-4. (Continued)

```

0018 320800      STA      Selected$Disk      ;Save selected disk number
                                           ;Set up to return DPH address
001B 6F          MOV      L,A                ;Make disk into word value
001C 2600        MVI      H,0
                                           ;Compute offset down disk parameter
                                           ; header table by multiplying by
                                           ; parameter header length (16 bytes)
                                           ;*2
001E 29          DAD      H                ;*4
001F 29          DAD      H                ;*8
0020 29          DAD      H                ;*16
0021 29          DAD      H                ;Get base address
0022 110F00     LXI      D,Disk$Parameter$Headers
0025 19          DAD      D                ;DE -> appropriate DPH
0026 E5          PUSH     H                ;Save DPH address
                                           ;Access disk parameter block in order
                                           ; to extract special prefix byte that
                                           ; identifies disk type and whether
                                           ; deblocking is required
                                           ;
0027 110A00     LXI      D,10                ;Get DPB pointer offset in DPH
002A 19          DAD      D                ;DE -> DPB address in DPH
002B 5E          MOV      E,M                ;Get DPB address in DE
002C 23          INX      H
002D 56          MOV      D,M                ;DE -> DPB
002E EB          XCHG

SELDSK$Set$Disk$Type:
002F 2B          DCX      H                ;DE -> prefix byte
0030 7E          MOV      A,M                ;Get prefix byte
0031 E670        ANI      Disk$Type$Mask        ;Isolate disk type
0033 320900     STA      Disk$Type            ;Save for use in low-level driver
0036 7E          MOV      A,M                ;Get another copy of prefix byte
0037 E680        ANI      Need$Deblocking      ;Isolate deblocking flag
0039 320A00     STA      Deblocking$Required ;Save for use in low-level driver
                                           ;Additional code to check if spare
                                           ; directory for given disk has already
                                           ; been read in.

003C 7E          MOV      A,M                ;Get physical disk number
003D E60F        ANI      Physical$Disk$Mask
003F 320B00     STA      Selected$Physical$Disk ;Save for low-level drivers

0042 5F          MOV      E,A                ;Make into word
0043 1600        MVI      D,0
0045 210400     LXI      H,Spare$Dir$In$Memory ;Make pointer into table
0048 19          DAD      D

0049 7E          MOV      A,M                ;Get flag
004A B7          ORA      A
004B C27700     JNZ      Dir$In$Memory        ;Spare directory already in memory
004E 34          INR      M                ;Set flag

004F 210000     LXI      H,Spare$Directories ;Create pointer to spare
0052 19          DAD      D                ; spare directory (added twice)
0053 19          DAD      D                ; as table has word entries)
                                           ;HL -> word containing directory addr.

0054 5E          MOV      E,M                ;Spare directory address in DE
0055 23          INX      H                ;HL -> spare directory
0056 56          MOV      D,M
0057 EB          XCHG

0058 220600     SHLD     Selected$Spare$Directory ;Save for use in physical
                                           ; drivers later on

005B 110000     LXI      D,Spare$Track        ;Track containing spare directory
005E 3A0B00     LDA      Selected$Physical$Disk
0061 47          MOV      B,A
0062 3E04        MVI      A,Spare$Sector      ;Sector containing spare directory
0064 0E18        MVI      C,Spare$Length/8   ;Number of bytes in spare directory / 8
0066 CDB500     CALL     Absolute$Read      ;Read in spare directory - without
                                           ; using bad sector management

```

Figure 9-4. (Continued)

```

0069 2A0600      LHLD   Selected$Spare$Directory ;Set end marker
006C 11C000      LXI    D,Spare$Length           ; at back end of spare directory
006F 19          DAD    D
0070 110080      LXI    D,Maximum$Track         ;Use maximum track number
0073 73          MOV    M,E
0074 23          INX   H
0075 3602        MVI    M,D

Dir$In$Memory:
0077 E1          POP    H                       ;Recover DPH pointer
0078 C9          RET

;
;
;   In the low-level disk drivers, the following code must be
;   inserted just before the disk controller is activated to
;   execute a read or a write command.
;
0079 2A0C00      LHLD   Disk$Track              ;Get track number from disk
;                               ; controller command table
007C EB          XCHG                      ;DE = track
007D 2A0600      LHLD   Selected$Spare$Directory ;HL -> spare directory
0080 2B          DCX    H                       ;Back up one entry
0081 2B          DCX    H                       ; (3 bytes)
0082 2B          DCX    H

0083 3A0E00      LDA    Disk$Sector            ;Get sector number
0086 4F          MOV    C,A                       ;Save for later

0087 06FF        MVI    B,OFFH                  ;Set counter (biased -1)

Check$Next$Entry:
0089 23          INX    H                       ;Update to next (or first) entry
Check$Next$Entry1:
008A 23          INX    H
Check$Next$Entry2:
008B 23          INX    H

008C 04          INR    B                       ;Update count

IF      Sector$Sparing

;If sparing is used, the
; end of the table is indicated
; by an entry with the track number
; = to maximum track number
LXI    D,Maximum$Track         ;Get maximum track number
CALL   CMPM                    ;Compare DE to (HL), (HL+1)
JZ     Not$Bad$Sector          ;End of table reached

ENDIF

;Note: For sector skipping
; the following search loop will
; terminate when the requested track
; is less than that in the table.
;This will always happen when the
; maximum track number is encountered
; at the end of the table.

008D EB          XCHG                      ;DE -> table entry
008E 2A0C00      LHLD   Disk$Track              ;Get requested track
0091 EB          XCHG                      ;DE = req. track, HL -> table entry
0092 CDCD00      CALL   CMPH                    ;Compare req. track to table entry

IF      Sector$Sparing

;Use the following code for
; sector sparing
JNZ    Check$Next$Entry        ;Track does not match
INX    H                       ;HL -> MS byte of track
INX    H                       ;HL -> sector
MOV    A,C                     ;Get requested sector
CMP    M                       ;Compare to table entry
JNZ    Check$Next$Entry2       ;Sector does not match

;Track and sector match, so
; substitute spare track and
; appropriate sector

```

Figure 9-4. (Continued)

```

LXI H,Spare$Track ;Get track number used for spare
; sectors
SHLD Disk$Track ;Substitute track

MVI A,First$Spare$Sector ;Get first sector number
ADD B ;Add on matched directory
; entry number
STA Disk$Sector ;Substitute sector
ENDIF

IF Sector$Skipping ;Use the following code for
; sector skipping
;The object is to find the
; entry in the table which
; is greater or equal to the
; requested sector/track

0095 CA9E00 JZ Tracks$Match ;Possible match of track and sector
0098 D2AC00 JNC Compute$Increment ;Requested track < table entry
009B C3B900 JMP Check$Next$Entry ;Requested track > table entry

Tracks$Match:
009E 23 INX H ;HL -> MS byte of track
009F 23 INX H ;HL -> sector
00A0 77 MOV M,A ;Get sector from table

00A1 B9 CMP C ;Compare with requested sector
00A2 CAAB00 JZ Sectors$Match ;Track/sector matches
00A5 D2AC00 JNC Compute$Increment ;Req. trk/sec < spare trk/sec
00A8 C3BB00 JMP Check$Next$Entry2 ;Move to next table entry

Sectors$Match:
00AB 04 INR B ;If track and sectors match with
; a table entry, then an additional
; sector must be skipped

Compute$Increment:
;B contains number of cumulative
; number of sectors to skip
00AC 79 MOV A,C ;Get requested sector
00AD 80 ADD B ;Skip required number
00AE 0612 MVI B,Sectors$Per$Track ;Determine final sector number
; and track increment
00B0 CDC300 CALL DIV$A$BY$B ;Returns C = quotient, A = remainder
00B3 320E00 STA Disk$Sector ;A = new sector number

00B6 59 MOV E,C ;Make track increment a word
00B7 1600 MVI D,0
00B9 2A0C00 LHL Disk$Track ;Get requested track
00BC 19 DAD D ;Add on increment
00BD 220C00 SHLD Disk$Track ;Save updated track
ENDIF

Not$Bad$Sector:
;Either track/sector were not bad,
; or requested track and sector have
; been updated.
00C0 C3D500 JMP Read$Write$Disk ;Go to physical disk read/write
;
; IF Sector$Skipping ;Subroutine required for skipping
; routine
;
;
; DIV$A$BY$B
; Divide A by B
;
;
; This routine divides A by B, returning the quotient in C
; and the remainder in A.
;
;
; Entry parameters
;
; A = dividend
; B = divisor
;
; Exit parameters
;

```

Figure 9-4. (Continued)


```

;
;           A = remainder
;           C = quotient
;
;
; DIV#A#BY#B:
00C3 OE00      MVI     'C,0           ;Initialize quotient
;
; DIV#A#BY#B$Loop:
00C5 0C       INR     C           ;Increment quotient
00C6 90       SUB     B           ;Subtract divisor
00C7 F2C500   JP      DIV#A#BY#B$Loop ;Repeat if result still +ve
00CA 0D       DCR     C           ;Correct quotient
00CB 80       ADD     B           ;Correct remainder
00CC C9       RET
;
;           ENDF
;
;
; CMPM
; Compare memory
;
;
; This subroutine compares the contents of DE to (HL) and (HL+1)
; returning with the flags as though the subtraction (HL) - DE
; were performed.
;
; Entry parameters
;
;           HL -> word in memory
;           DE = value to be compared
;
;
; Exit parameters
;
;           Flags set for (HL) - DE
;
;
; CMPM:
00CD 7E       MOV     A,M           ;Get MS byte
00CE BA       CMP     D
00CF C0       RNZ           ;Return now if MS bytes unequal
00D0 23       INX     H           ;HL -> LS byte
00D1 7E       MOV     A,M           ;Get LS byte
00D2 BB       CMP     E
00D3 2B       DCX     H           ;Return with HL unchanged
00D4 C9       RET
;
;
; Absolute$Read:
;           ;The absolute read (and write) routines
;           ; access the specified sector and track
;           ; without using bad sector management.
;
; Entry parameters
;
;           HL -> Buffer
;           DE = Track
;           A = Sector
;           B = Physical disk drive number
;           C = Number of bytes to read / 8
;
; Set up disk controller command block with parameters in
; registers, then initiate read operation by falling through
; into Read$Write$Disk code below.
;
;
; Read$Write$Disk:
;
; -----
; The remainder of the low level disk drivers follow,
; reading the required sector and track.
; -----
;
;
; Spare directory declarations
;
;
; Note: The disk format utility creates an initial spare
; directory with track/sector entries for those track/sectors
; that it finds are bad. It fills the remainder of the
; directory with OFFH's (these serve to terminate the
; searching of the directory).

```

Figure 9-4. (Continued)

```

;
;
00C0 = Spare$Length EQU 64 * 3 ;64 Entries, 3 bytes each
; Byte 0,1 = track
; Byte 2 = sector

Spare$Directory$0:
00D5 DS Spare$Length ;Spare directory itself
0195 DS 2 ;Set to maximum track number by SELDSK as
; a safety precaution. The FORMAT utility
; puts the maximum track number into all
; unused entries in the spare directory.

Spare$Directory$1:
0197 DS Spare$Length ;Spare directory itself
0257 DS 2 ;End marker

```

Figure 9-4. Bad sector management (continued)

Improving Error Messages

The final extension to BIOS error handling discussed here is in disk-driver error-message handling. The subroutine shown in the example BIOS in Figure 8-10, although a significant improvement on the messages normally output by the BDOS, did not advise the user of the most suitable course of action for each error. Figure 9-5 shows an improved version of the error message processor.

```

; This shows slightly more user-friendly error processor
; for disk errors than that shown in the enhanced BIOS
; in Figure 8-10.
; This version outputs a recommended course of action
; depending on the nature of the error detected.
; Code that remains unchanged from Figure 8-10 has been
; abbreviated.
;
; Dummy equates and data declarations needed to get
; an error free assembly of this example.
;
0001 = Floppy$Read$Code EQU 01H ;Read command for controller
0002 = Floppy$Write$Code EQU 02H ;Write command for controller
;
0000 00 Disk$Hung$Flag: DB 0 ;Set NZ when watchdog timer times
; out
0258 = Disk$Timer EQU 600 ;10-second delay (16.66ms tick)
;
0043 = Disk$Status$Block EQU 43H ;Address in memory where controller
; returns status
;Values from controller command table

0001 00 Floppy$Command: DB 0
0002 00 Floppy$Head: DB 0
0003 00 Floppy$Track: DB 0
0004 00 Floppy$Sector: DB 0

```

Figure 9-5. User-friendly disk-error processor

```

0005 00.    Deblocking$Required:  DB    0    ;Flag set by SELDSK according
; to selected disk type

0006 00    Disk$error$Flag:      DB    0    ;Error flag returned to BDOS
;

0007 00    In$Buffer$Disk:      DB    0    ;Logical disk Id. relating to current
; disk sector in deblocking buffer
;
; Equates for Messages
;
0007 =     BELL    EQU    07H    ;Sound terminal bell
000D =     CR      EQU    0DH    ;Carriage return
000A =     LF      EQU    0AH    ;Line feed
;
0005 =     BDOS    EQU    5      ;BDOS entry point (for system reset)
;
;
;
No$Deblock$Retry:
;-----
; Omitted code to set up disk controller command table
; and initiate the disk operation
;-----
0008 C31500    JMP      Wait$For$Disk$Complete
;
;
Write$Physical:
;Write contents of disk buffer to
; correct sector
000B 3E02     MVI    A,Floppy$Write$Code ;Get write function code
000D C31200    JMP      Common$Physical ;Go to common code
Read$Physical:
;Read previously selected sector
; into disk buffer
0010 3E01     MVI    A,Floppy$Read$Code ;Get read function code
Common$Physical:
0012 320100    STA    Floppy$Command ;Set command table
;
Deblock$Retry:
;Re-entry point to retry after error
;-----
; Omitted code sets up disk controller command block
; and initiates the disk operation
;-----
;
Wait$For$Disk$Complete:
;Wait until disk status block indicates
; operation has completed, then check
; if any errors occurred
;On entry HL -> disk control byte
;Ensure hung flag clear
0015 AF      XRA    A
0016 320000    STA    Disk$Hung$Flag

0019 213100    LXI    H,Disk$Timed$Out ;Set up watchdog timer
001C 015802    LXI    B,Disk$Timer ;Time delay
001F CD3B03    CALL   Set$Watchdog

Disk$Wait$Loop:
0022 7E      MOV    A,M ;Get control byte
0023 B7      ORA    A
0024 CA3700    JZ     Disk$Complete ;Operation done

0027 3A0000    LDA    Disk$Hung$Flag ;Also check if timed out
002A B7      ORA    A
002B C29F02    JNZ    Disk$error ;Will be set to 40H

002E C32200    JMP    Disk$Wait$Loop

Disk$Timed$Out:
;Control arrives here from watchdog
; routine itself -- so this is effectively
; part of the interrupt service routine.
0031 3E40     MVI    A,40H ;Set disk hung error code
0033 320000    STA    Disk$Hung$Flag ; into error flag to pull
; control out of loop
0036 C9      RET ;Return to watchdog routine

```

Figure 9-5. (Continued)

```

0037 010000      Disk$Complete:
                LXI      B,0                ;Reset watchdog timer
                                                ;HL is irrelevant here
003A CD3B03          CALL      Set$Watchdog

003D 3A4300          LDA      Disk$Status$Block      ;Complete -- now check status
0040 FE80            CPI      80H                ;Check if any errors occurred
0042 DA9F02          JC       Disk$Error          ;Yes

;
; Disk$Error$Ignore:
0045 AF             XRA      A                    ;No
0046 320600          STA      Disk$Error$Flag      ;Clear error flag
0049 C9             RET

;
; Disk error message handling
;
;
; Disk$Error$Messages:
; This table is scanned, comparing the
; disk error status with those in the
; table. Given a match, or even when
; the end of the table is reached, the
; address following the status value
; points to the correct advisory message text.
; Following this is the address of an
; error description message.

004A 40             DB      40H
004B B0019500        DW      Disk$Advice1,Disk$Msg$40
004F 41             DB      41H
0050 C9019A00        DW      Disk$Advice2,Disk$Msg$41
0054 42             DB      42H
0055 E301A400        DW      Disk$Advice3,Disk$Msg$42
0059 21             DB      21H
005A 0702B400        DW      Disk$Advice4,Disk$Msg$21
005E 22             DB      22H
005F 1B02B900        DW      Disk$Advice5,Disk$Msg$22
0063 23             DB      23H
0064 1B02C000        DW      Disk$Advice5,Disk$Msg$23
0068 24             DB      24H
0069 3D02D200        DW      Disk$Advice6,Disk$Msg$24
006D 25             DB      25H
006E 3D02DE00        DW      Disk$Advice6,Disk$Msg$25
0072 11             DB      11H
0073 5302F100        DW      Disk$Advice7,Disk$Msg$11
0077 12             DB      12H
0078 5302FF00        DW      Disk$Advice7,Disk$Msg$12
007C 13             DB      13H
007D 53020C01        DW      Disk$Advice7,Disk$Msg$13
0081 14             DB      14H
0082 53021A01        DW      Disk$Advice7,Disk$Msg$14
0086 15             DB      15H
0087 53022901        DW      Disk$Advice7,Disk$Msg$15
008B 16             DB      16H
008C 53023501        DW      Disk$Advice7,Disk$Msg$16
0090 00             DB      0                    ;<= Terminator
0091 53024501        DW      Disk$Advice7,Disk$Msg$Unknown ;Unmatched code

0005 =             DEM$Entry$Size EQU      5          ;Entry size in error message table
;
;
; Message texts
;
0095 48756E6700Disk$Msg$40: DB      'Hung',0          ;Timeout message
009A 4E6F742052Disk$Msg$41: DB      'Not Ready',0
00A4 5772697465Disk$Msg$42: DB      'Write Protected',0
00B4 4461746100Disk$Msg$21: DB      'Data',0
00B9 466F726D61Disk$Msg$22: DB      'Format',0
00C0 4D69737369Disk$Msg$23: DB      'Missing Data Mark',0
00D2 4275732054Disk$Msg$24: DB      'Bus Timeout',0
00DE 436F6E7472Disk$Msg$25: DB      'Controller Timeout',0
00F1 4472697665Disk$Msg$11: DB      'Drive Address',0
00FF 4865616420Disk$Msg$12: DB      'Head Address',0
010C 547261636BDisk$Msg$13: DB      'Track Address',0

```

Figure 9-5. (Continued)

```

011A 536563746FDisk$Msg$14:    DB      'Sector Address',0
0129 4275732041Disk$Msg$15:    DB      'Bus Address',0
0135 496C6C6567Disk$Msg$16:    DB      'Illegal Command',0
0145 556E6B6E6FDisk$Msg$Unknown: DB      'Unknown',0
;
;Disk$EM$1:
014D 070DOA                    DB      BELL,CR,LF ;Main disk error message -- part 1
0150 4469736B20                DB      'Disk ',0
;
;Error text output next
;
;Disk$EM$2:
0156 204572726F                DB      ' Error ('
015E 0000 Disk$EM$Status: DB      0,0 ;Status code in hex
0160 290DOA2020                DB      ')',CR,LF,' Drive '
016E 00 Disk$EM$Drive: DB      0 ;Disk drive code, A,B...
016F 2C20486561                DB      ', Head '
0176 00 Disk$EM$Head: DB      0 ;Head number
0177 2C20547261                DB      ', Track '
017F 0000 Disk$EM$Track: DB      0,0 ;Track number
0181 2C20536563                DB      ', Sector '
018A 0000 Disk$EM$Sector: DB      0,0 ;Sector number
018C 2C204F7065                DB      ', Operation - '
019A 00 DB 0 ;Terminator
;
019B 526561642EDisk$EM$Read: DB 'Read.',0 ;Operation names
01A1 5772697465Disk$EM$Write: DB 'Write.',0
;
01A8 0D0A202020Disk$Advice0: DB CR,LF,' ',0
01B0 436865636BDisk$Advice1: DB 'Check disk loaded, Retry',0
01C9 506F737369Disk$Advice2: DB 'Possible hardware problem',0
01E3 5772697465Disk$Advice3: DB 'Write enable if correct disk, Retry',0
0207 5265747279Disk$Advice4: DB 'Retry several times',0
021B 5265666F72Disk$Advice5: DB 'Reformat disk or use another disk',0
023D 4861726477Disk$Advice6: DB 'Hardware error, Retry',0
0253 4861726477Disk$Advice7: DB 'Hardware or Software error, Retry',0
;
0275 2C206F7220Disk$Advice9: DB ', or call for help if error persists',CR,LF
;
;Disk>Action$Confirm:
029B 00 DB 0 ;Set to character entered by user
029C 0D0A00 DB CR,LF,0
;
; Disk error processor
;
; This routine builds and outputs an error message.
; The user is then given the opportunity to:
;
; R -- retry the operation that caused the error
; I -- ignore the error and attempt to continue
; A -- abort the program and return to CP/M
;
;Disk$Error:
029F F5 PUSH PSW ;Preserve error code from controller
02A0 215E01 LXI H,Disk$EM$Status ;Convert code for message
02A3 CD3B03 CALL CAH ;Converts A to hex
;
02A6 3A0700 LDA In$Buffer$Disk ;Convert disk id. for message
02A9 C641 ADI 'A'
02AB 326E01 STA Disk$EM$Drive ;Make into letter
;
02AE 3A0200 LDA Floppy$Head ;Convert head number
02B1 C630 ADI '0'
02B3 327601 STA Disk$EM$Head
;
02B6 3A0300 LDA Floppy$Track ;Convert track number
02B9 217F01 LXI H,Disk$EM$Track
02BC CD3B03 CALL CAH
;
02BF 3A0400 LDA Floppy$Sector ;Convert sector number
02C2 218A01 LXI H,Disk$EM$Sector
02C5 CD3B03 CALL CAH
;
02C8 214D01 LXI H,Disk$EM$1 ;Output first part of message
02CB CD3B03 CALL Output$error$message
    
```

Figure 9-5. (Continued)

```

02CE F1          POP      PSW          ;Recover error status code
02CF 47          MOV      B,A           ;For comparisons
02D0 214500     LXI      H,Disk$Error$Messages - DEM$Entry$Size
                                ;HL -> table -- one entry
                                ;For loop below
02D3 110500     LXI      D,DEM$Entry$Size
                                ;HL -> table -- one entry
                                ;For loop below
02D6 19          DAD      D           ;Move to next (or first) entry

02D7 7E          MOV      A,M           ;Get code number from table
02D8 B7          ORA      A           ;Check if end of table
02D9 CAE302     JZ       Disk$Error$Matched ;Yes, pretend a match occurred
02DC B8          CMP      B           ;Compare to actual code
02DD CAE302     JZ       Disk$Error$Matched ;Yes, exit from loop
02E0 C3D602     JMP      Disk$Error$Next$Code ;Check next code

;
Disk$Error$Matched:
02E3 23          INX      H           ;HL -> advisory text address
02E4 5E          MOV      E,M
02E5 23          INX      H
02E6 56          MOV      D,M         ;DE -> advisory test
02E7 D5          PUSH   D           ;Save for later

02E8 23          INX      H           ;HL -> message text address
02E9 5E          MOV      E,M         ;Get address into DE
02EA 23          INX      H
02EB 56          MOV      D,M

02EC EB          XCHG                    ;HL -> text
02ED CD3B03     CALL   Output$error$Message ;Display explanatory text

02F0 215601     LXI      H,Disk$EM$2    ;Display second part of message
02F3 CD3B03     CALL   Output$error$Message

02F6 219B01     LXI      H,Disk$EM$Read ;Choose operation text
                                ; (assume a read)
02F9 3A0100     LDA      Floppy$Command ;Get controller command
02FC FE01     CPI      Floppy$Read$Code
02FE CA0403     JZ       Disk$Error$Read ;Yes
0301 21A101     LXI      H,Disk$EM$Write ;No, change address in HL
                                ;
Disk$Error$Read:
0304 CD3B03     CALL   Output$error$Message ;Display operation type
0307 21A801     LXI      H,Disk$Advice0 ;Display leading blanks
030A CD3B03     CALL   Output$error$Message

030D E1          POP      H           ;Recover advisory text pointer
030E CD3B03     CALL   Output$error$Message

0311 217502     LXI      H,Disk$Advice9 ;Display trailing component
0314 CD3B03     CALL   Output$error$Message

;
Disk$Error$Request$Action:
0317 CD3B03     CALL   Request$User$Choice ;Ask the user what to do next
                                ;Display prompt and get single
                                ; character response (folded to
                                ; uppercase)
                                ;Retry
031A FE52     CPI      'R'         ;Retry
031C CA2C03     JZ       Disk$Error$Retry
031F FE41     CPI      'A'         ;Abort?
0321 CA3603     JZ       System$Reset
0324 FE49     CPI      'I'         ;Ignore?
0326 CA4500     JZ       Disk$Error$Ignore
0329 C31703     JMP      Disk$Error$Request$Action

;
Disk$Error$Retry:
                                ;The decision on where to return to
                                ; depends on whether the operation
                                ; failed on a deblocked or
                                ; nondeblocked drive
032C 3A0500     LDA      Deblocking$Required
032F B7          ORA      A
0330 C21500     JNZ      Deblock$Retry
0333 C30800     JMP      No$Deblock$Retry

```

Figure 9-5. (Continued)

```

;
; System$Reset:                                ;This is a radical approach, but
0336 0E00          MVI    C,0                    ; it does cause CP/M to restart
0338 CD0500        CALL   BD0S                    ;System reset
;
;          Omitted subroutines (listed in full in Figure 8-10)
;
; Set$Watchdog:                                ;Set watchdog timer (to number of "ticks" in BC, and
;          ; to transfer control to (HL) if timer hits zero).
CAH:              ;Convert A to two ASCII hex characters, storing
;          ; the output in (HL) and (HL+1)
Output$error$Message: ;Display the 00-byte terminated error message
;          ; pointed to by HL. Output is directed only to
;          ; those console devices not being used for list
;          ; output as well.
Request$User$Choice: ;Display prompt "Enter R, A, I..." and return
033B C9           RET                          ; single keyboard character (uppercase) in A
;          ;Dummy

```

Figure 9-5. User-friendly disk-error processor (continued)

10

Basic Debugging Techniques
Debug Subroutines
Software Tools for Debugging
Bringing Up CP/M for the First Time
Debugging the CP/M Bootstrap
Loader
Debugging the BIOS
Live Testing a New BIOS

Debugging A New CP/M System

This chapter deals with some of the problems you will face bringing up CP/M on a computer system for the first time or enhancing it once it is up and running on your system.

In the first case, when CP/M does not yet run on your computer, you may be writing the complete BIOS yourself, although you can model what you do on the example BIOS provided on the CP/M release diskette and the example code from Chapter 6.

In the second case, you can extend the existing BIOS by adding code—from the examples in Chapters 8 and 9, code from computer magazines, or code you create yourself. To do this, you will need access to the BIOS source code—a problem if the manufacturer of your computer does not make it available. In general, however, the BIOS source code is included with the system or can be obtained at nominal or no cost. If you cannot obtain the source code, you can, of

course, take the bull by the horns and reimplement CP/M on your system. This may require many hours of disassembling the current BIOS machine code to find out how to access all the various ports and how to control the devices to which they are connected.

Although the BIOS is the major component of a new CP/M implementation, remember that it is only the beginning—you can spend the same amount of time and effort getting the bootstrap loader and all the utilities to function.

Basic Debugging Techniques

Before getting involved in the details of how to debug a CP/M implementation, it is worth considering the nature of the task. Some quotations that are appropriate here:

“Program testing can be used to show the presence of bugs, but never to show their absence.” —Dijkstra

“We call them bugs because to call them mistakes would be psychologically unacceptable.” —Hopkins

“Constants aren't, variables won't.” —Osborne

Debugging is the name we give to the process of executing programs and ascertaining whether the programs are running correctly. “Correctly” means in accordance with the mental model we have built of how the program should behave, subject to the constraints imposed by the physical hardware. Therein lies the first of the problems; you and the hardware are the arbiters of correct performance. The hardware is usually unforgiving; if there is a flaw in the way you program it, it will either be dramatically “uncooperative” or not work at all. As for how you perceive the system, several fairly simple tests, along with attempts to use the system for useful work for a few days, will shake the system down fairly well. The most difficult problems will be with intermittent failures or logical contradictions.

Computers are deterministic. That is, if you start from a known state and perform a known series of operations, the computer will always yield the same results. To achieve a known state is not so difficult—resetting the system and clearing memory will do it. Performing a known series of operations just means running the program again, although if you are using interrupts, you cannot truthfully say that exactly the same operations are being performed, because the interrupts will not happen at *exactly* the same time as before.

The “Orville Wright” Approach

Your role in debugging a new CP/M system is comparable to the popular, though untrue, idea of the way the Wright brothers developed flying machines:

build a machine, take it to the top of a hill, throw it off, and, when it crashes, examine the debris to discover what went wrong.

Each time you do an assembly and test, you are building the aircraft and lobbing it off the edge of a cliff. Each time it crashes, you examine the wreckage and try to determine the possible cause.

This is a highly inferential process. With the wreckage as a starting point, you use inference and intuition to extrapolate the real problem and the correction for it.

Built-In Debug Code

The single most important concept that you will need in testing CP/M systems is the same as that used in the modern day “black box” flight recorder. This device is essentially a multi-channel tape recorder that records all of the relevant conditions of the aircraft, its height, altitude, throttle settings, flap settings, and even the voice communications among crew members. If the airplane crashes, investigators can replay the information and understand what happened during the flight.

Applying this concept to debugging CP/M means that you must build into your code some method for recording what it is doing, so that if the system crashes, you can see what it was doing. Make the code tell you what went wrong.

The debug code should be designed at the same time as the rest of the program. Plan the debugging code while the design is still on the drawing board. The source code for debugging should be a permanent part of the BIOS. Use conditional assembly to “IF” out most of the debug code from the final version, or make the code sensitive to a flag in the configuration block so that you can re-enable the debug code at a moment’s notice if the system begins to behave strangely.

The more meaningful the debug output data, the less you will have to guess at what is wrong, and therefore the less painful and time-consuming the debugging process will be. Make the output intelligible to others who may use it or yourself several months hence. Data that tells you what is happening is more useful than internal hexadecimal values, particularly if someone else must interpret it or relay it to you over the telephone.

Debug Subroutines

Many programmers do their debugging on a casual “catch as catch can” basis because they are overwhelmed by the task of building the necessary tools. Others are too eager to start on a new program to take a few extra hours or days to build debug subroutines.

To help solve this problem, the following section provides some ready-made debugging tools that can be used “as is.” Each of these routines has been thor-

oughly debugged (there's nothing worse than debug code with bugs in it!) and has been used in actual program testing.

Overall Design Philosophy

Some common methods run through the examples that follow. These include displaying meaningful "captions" (including the specific address that called the debug routine), grouping all debugging code together, preserving the contents of all registers, and setting up the stack area in a standard way.

Debug Code Captions When the contents of registers or memory are output as part of a debugging process, a caption of explanatory text describing the values should be displayed. For example, rather than displaying the contents of the A register like this,

```
A = 1F
```

you can use a meaningful caption such as:

```
Transaction Code A = 1F.
```

When you write additional debugging code, especially if you need to add it to an existing routine, it is cumbersome to have to write the call to the debug routine and then search through the source code to find a convenient place to put an ASCII caption string. A caption string several pages removed from the point where it is referenced makes for problems when you want to relate the debug output on the screen or listing to the source code itself. Therefore, all of the routines that follow allow you to declare the caption strings "in-line" like this:

```
IF          DEBUG
CALL       Debug$Routine
DB        'Caption string here',CR,LF,0
ENDIF

MVI       .....           ;Next instruction
```

All of the following routines that output a caption recognize one specific 8-bit value in the caption string. If they encounter a value of 0ADH (mnemonic for ADdress), they will output the address of the byte following the call to the debug routine. For example,

```
0210      CALL       Debug$Routine
0213      DB        0ADH,'Caption string',0
```

will cause the routine to display the following:

```
0213 Caption string
```

This identifies the point in your program from which the debug routine was called, and thus avoids any possible ambiguity between different calls to the same debug routine with similar captions.

Grouping Debug Code Grouping all the debug code together lends itself to using conditional assembly with IF/ENDIF statements.

Setting Up the Stack Area All of the following routines preserve the CPU registers so that there are no side effects from using them. All of them assume that they can use the stack pointer and that there is sufficient room in the stack area. Hence you will need to declare adequate stack space for your main code and for the debug routines. Fill the stack area with a known pattern like this:

```

        DW    9999H, 9999H, 9999H, 9999H, 9999H, 9999H, 9999H, 9999H, 9999H
        DW    9999H, 9999H, 9999H, 9999H, 9999H, 9999H, 9999H, 9999H, 9999H
        DW    9999H, 9999H, 9999H, 9999H, 9999H, 9999H, 9999H, 9999H, 9999H
Stack$Area:      ;Label the upper end of the area

```

Then, during debugging, you can examine the stack area and determine how much of it is unused. For example, if you looked at the stack area you might see something like this:

```

                                "Low-water mark"
                                V
99 99 99 99 99 99 99 99 99 99 99 99 09 15 43 42
01 29 00 00 1A 2B 10 FF FF 39 02 ED 11 01 37 44
DD 00 00 11 1A 23 31 00 41 AE FE 00 01 10 70 C9

```

Stack area overflow can give arcane bugs; the program seems to leap off into space in a nondeterministic way. By setting up the stack area in this way, you can recognize an overflow condition easily.

Debug Initialization Before you can execute any of the debug subroutines in this chapter, you must make a call to the initialization subroutine, DB\$Init. The DB\$Init routine sets up some of the internal variables needed by the debug package. You may need to add some of your own initialization code here.

Console Output

Normally, you can use the CONOUT functions either via the BDOS (Function 2), or via the BIOS by calling the jump vector directly. You cannot do this when you need to debug console routines themselves, nor when you need to debug interrupt service routines. In the latter case, if an interrupt pulled control out of the CONOUT routine in the BIOS, you would get unwanted re-entrancy if the debug code again entered the CONOUT driver to display a caption. Therefore, the debug routines have been written to call their own local CONOUT routine, which is called DB\$CONOUT. DB\$CONOUT can be changed to call the BDOS, the BIOS, or a "private" polled output routine.

A counterpart DB\$CONIN routine for console input is provided for essentially the same reasons.

Controlling Debug Output

All output of debug routines in this chapter is controlled by a single master flag, `DB$Flag`. If this flag is nonzero, debug output will occur; if zero, all output is suppressed.

This flag can be set and cleared from any part of the program you are testing. It is especially useful when you need to debug a subroutine that is called many times from many different places. You can write additional code to enable debug output when certain conditions prevail; for example, when a particular track or sector is about to be written or when a character input buffer is almost full.

Two subroutines, `DB$On` and `DB$Off`, are shown that access the debug control flag. These, as their names suggest, turn debug output on and off.

Turning the debug output on and off from within the program can create a confusing display of debug output, lacking any apparent continuity. `DB$Off` gives you the option of outputting a character string indicating that debug output has been turned off.

Pass Counters

Another method of controlling debug output is to use a *pass counter*, enabling debug output only after control has passed through a particular point in the code a specific number of times.

Two subroutines are provided for this purpose. `DBSetPass` sets the pass counter to a specific value. `DB$Pass` decrements this pass count each time control is transferred to it. When the pass count hits zero, the debug control flag `DB$Flag` is nonzero and debug output begins.

Using pass counter techniques can save you time and effort in tracking down a problem that occurs only after the code has been running for several minutes.

Displaying Contents of Registers and Memory

Figure 10-2 shows a series of display subroutines, the primary one of which is `DB$Display`. It takes several parameters, depending on the information you want displayed. The generic call to `DB$Display` is as follows:

```
CALL      DB$Display
DB        Code      <- Indicates the data to be
                    displayed
{DW      Optional additional parameters}
DB        'Caption string',0
```

The codes that can be used in this call are shown in Table 10-1.

The only function that uses additional parameters is `DB$Memory`. This displays bytes from memory in hexadecimal and ASCII, using the start and finish

addresses following the call. Here is an example:

```
CALL    DB$Display
DB      DE$Memory
DW      Start$Address,End$Address
DB      'Caption string',0
```

Table 10-1. Codes for DB\$Display

Code	Value displayed
8-bit registers	
DB\$F	Condition Flags
DB\$A	Register A
DB\$B	Register B
DB\$C	Register C
DB\$D	Register D
DB\$E	Register E
DB\$H	Register H
DB\$L	Register L
Memory	
DB\$Memory	Bytes starting and ending at the addresses specified by the two word values following the code value.
16-bit registers	
DB\$BC	Register pair BC
DB\$DE	Register pair DE
DB\$HL	Register pair HL
DB\$SP	Stack Pointer
Byte values	
DB\$B\$BC	Byte addressed by BC
DB\$B\$DE	Byte addressed by DE
DB\$B\$HL	Byte addressed by HL
Word values	
DB\$W\$BC	Word addressed by BC
DB\$W\$DE	Word addressed by DE
DB\$W\$HL	Word addressed by HL

Debugging Program Logic

In addition to displaying the contents of registers and memory, you need to display the program's execution path, not in terms of addresses, but in terms of the *problem*. You can do this by displaying debug messages that indicate what decisions have been made by the program as it executes. For example, if your BIOS checks a particular value to see whether the system should read or write on a particular device, the debug routine should display a message like this:

```
Entering Disk Read Routine
```

This is more meaningful than just displaying the function code for the drivers—although you may want to display this as well, in case it has been set to some strange value.

Two subroutines are provided to display debug messages. They are DB\$MSG and DB\$MSGI. Both of these display text strings are terminated with a byte of 00H. You can see the difference between the two subroutines if you examine the way they are called.

DB\$MSG is called like this:

```
LXI H,Message$Text      ;HL -> text string
CALL DB$MSG
```

DB\$MSGI is called like this:

```
CALL DB$MSG
DB 0DH,0AH,'Message Text',0 ;In-line
```

DB\$MSGI is more convenient to use. If you decide that you need to add a message, you can declare the message immediately following the call. This also helps when you look at the listing, since you can see the complete text at a glance.

Use DB\$MSG when the text of the message needs to be selected from a table. Get the address of the text into HL and then call DB\$MSG to display it.

Creating Your Own Debug Displays

If you need to build your own special debug display routines, you may find it helpful to incorporate some of the small subroutines in the debug package. The following are the subroutines you may want to use:

DB\$CONOUT

Displays the character in the C register.

DB\$CONIN

Returns the next keyboard character in A.

DB\$CONINU

Returns the next keyboard character in A, converting lowercase letters to uppercase.

DB\$DHLH

Displays contents of HL in hexadecimal.

DB\$DAH

Displays contents of A in hexadecimal.

DB\$CAH

Converts contents of A to hexadecimal and stores in memory pointed at by HL.

DB\$Nibble\$To\$Hex

Converts the least significant four bits of A into an ASCII hexadecimal character in A.

DB\$CRLF

Displays a CARRIAGE RETURN/LINE FEED.

DB\$Colon

Displays the string “:”.

DB\$Blank

Displays a single space character.

DB\$Flag\$Save\$On

Saves the current state of the debug output control flag and then sets the flag “on” to enable debug output.

DB\$Flag\$Restore

Restores the debug output control flag to the state it was in when the *DB\$Flag\$Save\$On* routine was last called.

DB\$GHV

Gets a hexadecimal value from the keyboard, displaying a prompt message first. From one to four characters can be specified as the maximum number of characters to be input.

DB\$A\$To\$Upper

If the A register contains a lowercase letter, this converts it to an uppercase letter.

Debugging I/O Drivers

Debugging low-level device drivers creates special problems. The major one is that you do not normally want to read and write via actual hardware ports while you are debugging the code—either because doing so would cause strange things to happen to the hardware during the debugging, or because you are developing and debugging the drivers on a system different from the target hardware on which the drivers are to execute.

Before considering the solution, remember that the input and output instructions (IN and OUT) are each two bytes long. The first byte is the operation code

(0DBH for input, 0D3H for output), and the second byte is the port number to “input from” or “output to.”

Debug subroutines are provided here to intercept all IN and OUT instructions, displaying the port number and either accepting a hexadecimal value from the console and putting it into the A register (in the case of IN), or displaying the contents of the A register (for the OUT instruction).

IN and OUT instructions can be “trapped” by changing the operation code to one of two RST (restart) instructions. An RST is effectively a single-byte CALL instruction, calling down to a predetermined address in low memory. The debug routines arrange for JMP instructions in low memory to receive control when the correct RST is executed. The code that receives control can pick up the port number, display it, and then accept a hex value for the A register (for IN) or display the current contents of the A register (for OUT). The example subroutines shown later in this chapter use RST 4 in place of IN instructions, RST 5 for OUT.

Wherever you plan to use IN, use the following code:

```
IF          Debug
RST         4
ENDIF
IF          NOT Debug
DB         IN
ENDIF
DB         Port$Number
```

Note that you can use the IN operation code as the operand of a DB statement. The assembler substitutes the correct operation code.

Use the following code wherever you need to use an OUT instruction:

```
IF          Debug
RST         5
ENDIF
IF          NOT Debug
DB         OUT
ENDIF
DB         Port$Number
```

When the RST 4 (IN) instruction is executed, the debug subroutine displays

```
1AB3 : Input from Port 01 : _
```

The “1AB3” is the address in memory of the byte containing the port number. It serves to pinpoint the IN instruction in memory. You can then enter one or two hexadecimal digits. These will be converted and put into the A register before control returns to the main program at the instruction following the byte containing the port number.

When the RST 5 (OUT) instruction is encountered, the debug subroutine displays

```
1AB5 : Output to Port 01 : FF
```

This identifies where the OUT instruction would normally be as well as the port number and the contents of the A register when the RST 5 (OUT) is executed.

Debugging Interrupt Service Routines

You can use a technique similar to that of the RST instruction just described to “fake” an interrupt. You preset the low-memory address for the RST instruction you have chosen for the jump into the interrupt service routine under test.

When the RST instruction is executed, control will be transferred into the interrupt service routine just as though an interrupt had occurred. You will need to intercept any IN or OUT instructions as described above—otherwise the code probably will go into an endless loop.

Before executing the RST instruction to fake the interrupt, load all the registers with known values. For example:

```
MVI      A,0AAH
LXI      B,0BCCCH
LXI      D,0DDEEH
LXI      H,01122H
RST      6           ;Fake interrupt
NOP
```

When control returns from the service routine, you can check to see that it restored all of the registers to their correct values. An interrupt service routine that does not restore all the registers can produce bugs that are very hard to find.

Check, too, that the stack pointer register has been restored and that the service routine did not require too many bytes on the stack.

You also can use the CALL instruction to transfer control to the interrupt service routine in order to fake an interrupt. RST and CALL achieve the same effect, but RST is closer to what happens when a real interrupt occurs. As it is a single-byte instruction, it also is easier to patch in.

Subroutine Listings

Figure 10-1 is a functional index to the source code listing for the debug subroutines shown in Figure 10-2. The listing’s commentary defines precisely how each debug subroutine is called.

Figure 10-3 shows the output from the debug testbed.

Software Tools for Debugging

In addition to building in debugging subroutines, you will need one of the following proprietary debug programs:

DDT (Dynamic Debugging Tool)

This program, included with the standard CP/M release, allows you to load programs, set and display memory and registers, trace through your program instruction by instruction, or execute it at full speed, but stopping

Start Line	Functional Component or Routines
00001	Debug subroutine's Testbed
00100	Test register display
00200	Test memory dump display
00300	Test register pair display
00400	Test byte indirect display
00500	Test DB\$On/Off
00600	Test DB\$Set\$Pass and DB\$Pass
00700	Test debug input/output
00800	Debug subroutines themselves
01100	DB\$Init - initialization
01200	DB\$CONINU - get uppercase keyboard character
01300	DB\$CONIN - get keyboard character
01400	DB\$CONOUT - display character in C
01500	DB\$On - enable debug output
01600	DB\$Off - disable debug output
01700	DB\$Set\$Pass - set pass counter
01800	DB\$Pass - execute pass point
01900	DB\$Display - main debug display routine
02200	Main display processing subroutines
02500	DB\$Display\$CALLA - display CALL's address
02600	DB\$DHLH - display HL in hexadecimal
02700	DB\$DAH - display A in hexadecimal
02800	DB\$CAH - convert A to hexadecimal in memory
02900	DB\$Nibble\$To\$Hex - convert LS 4 bits of A to hex.
02930	DB\$CRLF - display Carriage Return, Line Feed
02938	DB\$Colon - display ":"
02946	DB\$Blank - display ""
03100	DB\$MSGI - display in-line message
03147	DB\$MSG - display message addressed by HL
03300	DB\$Input - debug INput routine
03500	DB\$Output - debug OUTput routine
03700	DB\$Flag\$Save\$On - save debug flag and enable
03800	DB\$Flag\$Restore - restore debug control flag
03900	DB\$GHV - get hexadecimal value from keyboard
04100	DB\$A\$To\$Upper - convert A to upper case

Figure 10-1. Functional index for Figure 10-2

at certain addresses (called breakpoints). It also has a built-in mini-assembler and disassembler so you do not have to hand assemble any temporary code "patches" you add.

SID (Symbolic Interactive Debug)

Similar to DDT in many ways, SID has enhancements that are helpful if you use Digital Research's MAC (Macro Assembler) or RMAC (Relocating Macro Assembler). Both of these assemblers can be told to output a file

```

00001
00002
00003 ;
00004 ; Debug Subroutines
00005 ;
00006 ;<---- NOTE:
00007 ; The line numbers at the extreme left are included purely
00008 ; to reference the code from the text.
00009 ; There are deliberately induced discontinuities
00010 ; in the numbers in order to allow space for expansion.
00011 ;
00012 ; Because of the need to test these routines thoroughly,
00013 ; and in case you wish to make any changes, the testbed
00014 ; routine for the debug package itself has been left in
00015 ; in this figure.
00016 ;
00017 ; Debug testbed
00018 ;
00019 ;
00020 0100 ORG 100H
00021 START: LXI SP,Test$Stack ;Set up local stack
00022 0100 316B03 CALL DB$Init ;Initialize the debug package
00023 0103 CDEA04 CALL DB$On ;Enable debug output
00024 0106 CD1505 ;Simple test of A register display
00025 0109 3EAA MVI A,0AAH ;Preset a value in the A register
00026 010B 01CCBB LXI B,0BBCCCH ;Prefill all other registers, partly
00027 010E 11EEDD LXI D,0DDEEH ; to check the debug display, but
00028 0111 2111FF LXI H,0FF11H ; also to check register save/restore
00100 ;#
00101 ; Test register display
00102 ;
00103 0114 B7 ORA A ;Set M-flag, clear Z-flag, set E-flag
00104 0115 37 STC ;Set carry
00105 0116 CD5205 CALL DB$Display ;Call the debug routine
00106 0119 00 DB DB$F
00107 011A 466C616773 DB ^Flags',0
00108 ;
00109 0120 CD5205 CALL DB$Display ;Call the debug routine
00110 0123 02 DB DB$A
00111 0124 4120526567 DB ^A Register',0
00112 ;
00113 012F CD5205 CALL DB$Display ;Call the debug routine
00114 0132 04 DB DB$B
00115 0133 4220526567 DB ^B Register',0
00116 ;
00117 013E CD5205 CALL DB$Display ;Call the debug routine
00118 0141 06 DB DB$C
00119 0142 4320526567 DB ^C Register',0
00120 ;
00121 014D CD5205 CALL DB$Display ;Call the debug routine
00122 0150 08 DB DB$D
00123 0151 4420526567 DB ^D Register',0
00124 ;
00125 015C CD5205 CALL DB$Display ;Call the debug routine
00126 015F 0A DB DB$E
00127 0160 4520526567 DB ^E Register',0
00128 ;
00129 016B CD5205 CALL DB$Display ;Call the debug routine
00130 016E 0C DB DB$H
00131 016F 4820526567 DB ^H Register',0
00132 ;
00133 017A CD5205 CALL DB$Display ;Call the debug routine
00134 017D 0E DB DB$L
00135 017E 4C20526567 DB ^L Register',0
00200 ;#
00201 ; Test Memory Dump Display
00202 ;
00203 0189 CD5205 CALL DB$Display
00204 018C 18 DB DB$M ;Dump memory
00205 018D 08012801 DW 108H,128H ;Check start/end at nonmultiples
00206 0191 4D656D6F72 DB ^Memory Dump #1',0 ; of 10H
00207 ;
00208 01A0 CD5205 CALL DB$Display
00209 01A3 18 DB DB$M ;Dump memory
00210 01A4 00011F01 DW 100H,11FH ;Check start and end on displayed
00211 01A8 4D656D6F72 DB ^Memory Dump #2',0 ; line boundaries
00212 ;

```

Figure 10-2. Debug subroutines

```

00213 01B7 CD5205 CALL DB*Display
00214 01BA 18 DB DB*M ;Dump memory
00215 01BB 01010001 DW 101H,100H ;Check error handling where
00216 01BF 4D65D6F72 DB 'Memory Dump #3',0 ; start > end address
00217 ;
00218 01CE CD5205 CALL DB*Display
00219 01D1 18 DB DB*M ;Dump memory
00220 01D2 00010001 DW 100H,100H ;Check end-case of single byte
00221 01D6 4D65D6F72 DB 'Memory Dump #4',0 ; output
00300 ;#
00301 ;
00302 ; Test register pair display
00303 01E5 CD5205 CALL DB*Display ;Call the debug routine
00304 01E8 10 DB DB*BC
00305 01E9 4243205265 DB 'BC Register',0
00306 ;
00307 01F5 CD5205 CALL DB*Display ;Call the debug routine
00308 01F8 12 DB DB*DE
00309 01F9 4445205265 DB 'DE Register',0
00310 ;
00311 0205 CD5205 CALL DB*Display ;Call the debug routine
00312 0208 14 DB DB*HL
00313 0209 484C205265 DB 'HL Register',0
00314 ;
00315 0215 CD5205 CALL DB*Display ;Call the debug routine
00316 0218 16 DB DB*SP
00317 0219 5350205265 DB 'SP Register',0
00318 ;
00319 0225 013203 LXI B,Byte*BC ;Set up registers for byte tests
00320 0228 113303 LXI D,Byte*DE
00321 022B 213403 LXI H,Byte*HL
00400 ;#
00401 ;
00402 ; Test byte indirect display
00403 022E CD5205 CALL DB*Display ;Call the debug routine
00404 0231 1A DB DB*B*BC
00405 0232 4279746520 DB 'Byte at (BC)',0
00406 ;
00407 023F CD5205 CALL DB*Display ;Call the debug routine
00408 0242 1C DB DB*B*DE
00409 0243 4279746520 DB 'Byte at (DE)',0

00410 ;
00411 0250 CD5205 CALL DB*Display ;Call the debug routine
00412 0253 1E DB DB*B*HL
00413 0254 4279746520 DB 'Byte at (HL)',0
00414 ;
00415 0261 013503 LXI B,Word*BC ;Set up the registers for word tests
00416 0264 113703 LXI D,Word*DE
00417 0267 213903 LXI H,Word*HL
00418 ;
00419 026A CD5205 CALL DB*Display ;Call the debug routine
00420 026D 20 DB DB*W*BC
00421 026E 576F726420 DB 'Word at (BC)',0
00422 ;
00423 027B CD5205 CALL DB*Display ;Call the debug routine
00424 027E 22 DB DB*W*DE
00425 027F 576F726420 DB 'Word at (DE)',0
00426 ;
00427 028C CD5205 CALL DB*Display ;Call the debug routine
00428 028F 24 DB DB*W*HL
00429 0290 576F726420 DB 'Word at (HL)',0
00500 ;#
00501 ;
00502 ; Test DB*On/Off
00503 029D CD1D05 CALL DB*Off ;Disable debug output
00504 02A0 CDD607 CALL DB*MSGI ;Display in-line message
00505 02A3 0D0A546869 DB ODH,0AH,'This message should NOT appear',0
00506 ;
00507 02C4 CD1505 CALL DB*On
00508 02C7 CDD607 CALL DB*MSGI
00509 02CA 0D0A446562 DB ODH,0AH,'Debug output has been re-enabled.',0
00600 ;#
00601 ;
00602 ; Test pass count logic

```

Figure 10-2. (Continued)

```

00603 02EE CD1D05      CALL  DB#Off          ;Disable debug output
00604 02F1 CD2405      CALL  DB#Set#Pass    ;Set pass count
00605 02F4 1E00                ;
00606                ;
00607 02F6 3E22                MVI  A,34            ;Set loop counter greater than pass
00608                ; counter
00609                ;
00610                Test#Pass#Loop:
00610 02F8 CD3505      CALL  DB#Pass        ;Decrement pass count
00611 02FB CDD607      CALL  DB#MSGI       ;Display in-line message
00612 02FE 0D0A546869  DB    ODH,0AH,'This message should display 5 times',0
00613 0324 3D                DCR  A
00614 0325 C2F802      JNZ   Test#Pass$Loop
00700                ;#
00701                ; Test debug input/output
00702                ;
00703 0328 CD1D05      CALL  DB#Off          ;Check that debug IN/OUT
00704                ; must still occur when debug
00705                ; output is disabled.
00706 032B E7                RST   4              ;Debug input
00707 032C 11                DB    11H           ;Port number
00708 032D EF                RST   5              ;Debug output (value return from input)
00709 032E 22                DB    22H           ;Port number
00710                ;
00711 032F C30000      JMP   0              ;Warm boot at end of testbed
00712                ;
00713                ;
00714                ; Dummy values for byte and word displays
00715 0332 BC          Byte#BC:  DB    0BCH
00716 0333 DE          Byte#DE:  DB    0DEH
00717 0334 F1          Byte#HL:  DB    0F1H
00718                ;
00719 0335 0C0B        Word#BC:  DW    0B0CH
00720 0337 0E0D        Word#DE:  DW    0D0EH
00721 0339 010F        Word#HL:  DW    0F01H
00722                ;
00723 033B 9999999999   DW    9999H,9999H,9999H,9999H,9999H,9999H,9999H,9999H
00724 034B 9999999999   DW    9999H,9999H,9999H,9999H,9999H,9999H,9999H,9999H
00725 035B 9999999999   DW    9999H,9999H,9999H,9999H,9999H,9999H,9999H,9999H
00726                ;
00727                Test#Stack:
00728                ;
00729                ;
00730 0400                ORG   400H          ;To avoid unnecessary listings
00731                ; when only the testbed changes
00732                ;-----
00800                ;#
00801                ;
00802                ; Debug subroutines
00803                ;
00804                ;
00805                ; Equates for DB#Display codes
00806                ; These equates are the offsets down the table of addresses
00807                ; for various subroutines to be used.
00808                ;
00809 0000 =            DB#F   EQU    00      ;Flags
00810 0002 =            DB#A   EQU    02      ;A register
00811 0004 =            DB#B   EQU    04      ;B
00812 0006 =            DB#C   EQU    06      ;C
00813 0008 =            DB#D   EQU    08      ;D
00814 000A =           DB#E   EQU    10      ;E
00815 000C =           DB#H   EQU    12      ;H
00816 000E =           DB#L   EQU    14      ;L
00817 0010 =           DB#BC  EQU    16      ;BC
00818 0012 =           DB#DE  EQU    18      ;DE
00819 0014 =           DB#HL  EQU    20      ;HL
00820 0016 =           DB#SP  EQU    22      ;Stack pointer
00821 0018 =           DB#M   EQU    24      ;Memory
00822 001A =           DB#B#BC EQU    26      ;(BC)
00823 001C =           DB#B#DE EQU    28      ;(DE)
00824 001E =           DB#B#HL EQU    30      ;(HL)
00825 0020 =           DB#W#BC EQU    32      ;(BC+1),(BC)
00826 0022 =           DB#W#DE EQU    34      ;(DE+1),(DE)
00827 0024 =           DB#W#HL EQU    36      ;(HL+1),(HL)
00828                ;
00829                ;
00830                ; Equates
00831 0020 =            RST4   EQU    20H     ;Address for RST 4 - IN instruction

```

Figure 10-2. (Continued)

00832	0028 =	RST5	EQU	28H	;Address for RST 5 - OUT instruction
00833		;			
00834	0001 =	B#CONIN	EQU	1	;BDOS CONIN function code
00835	0002 =	B#CONOUT	EQU	2	;BDOS CONOUT function code
00836	000A =	B#READCONS	EQU	10	;BDOS read console function code
00837	0005 =	BDOS	EQU	5	;BDOS entry point
00838		;			
00839	0000 =	False	EQU	0	
00840	FFFF =	True	EQU	NOT False	
00841		;			
00842					;Equates to specify how DB#CONOUT
00843					; and DB#CONIN should perform
00844					; their input/output
00845	0000 =	DB#Polled\$I/O	EQU	False);
00846	0000 =	DB#BIOS\$I/O	EQU	False); Only one must be true
00847	FFFF =	DB#BDOS\$I/O	EQU	True);
00848		;			
00849					;Equates for polled I/O
00850	0001 =	DB#Status\$Port	EQU	01H	;Console status port
00851	0002 =	DB#Data\$Port	EQU	02H	;Console data port
00852		;			
00853	0002 =	DB#Input\$Ready	EQU	0000#0010B	;Incoming data ready
00854	0001 =	DB#Output\$Ready	EQU	0000#0001B	;Ready for output
00855		;			
00856					;Data for BIOS I/O
00857	0400 C3	BIOS#CONIN:	DB	JMP	;The initialization routine sets these
00858	0401 0000		DW	0	; two JMP addresses into the BIOS
00859	0403 C3	BIOS#CONOUT:	DB	JMP	
00860	0404 0000		DW	0	
00861		;			
00862		;			Main debug variables and constants
00863		;			
00864	0406 00	DB#Flag:	DB	0	;Main debug control flag
00865					; When this flag is nonzero, all debug
00866					output will be made. When zero, all
00867					debug output will be suppressed.
00868					; It is altered either directly by the user
00869					; or using the routines DB#On, DB#Off and
00870					; DB#Pass.
00871		;			
00872	0407 0000	DB#Pass\$Count:	DW	0	;Pass counter
00873					; When this is nonzero, calls to DB#Pass
00874					; decrement it by one. When it reaches
00875					; zero, the debug control flag, DB#Flag,
00876					; is set nonzero, thereby enabling
00877					; debug output.
00878		;			
00879		DB#Save\$HL:			;Save area for HL
00880	0409 00	DB#Save\$L:	DB	0	
00881	040A 00	DB#Save\$H:	DB	0	
00882					
00883	040B 0000	DB#Save\$SP:	DW	0	;Save area for stack pointer
00884	040D 0000	DB#Save\$RA:	DW	0	;Save area for return address
00885	040F 0000	DB#Call\$Address:	DW	0	;Starts out the same as DB#Save\$RA
00886					; but DB#Save\$RA gets updated during
00887					; debug processing. This value is
00888					; output ahead of the caption
00889		DB#Start\$Address:			;Start address for memory display
00890	0411 0000		DW	0	
00891		DB#End\$Address:			;End address for memory display
00892	0413 0000		DW	0	
00893		DB#Display\$Code:			;Display code requested
00894	0415 00		DB	0	
00895		;			
00896		;			
00897					;Stack area
00898	0416 9999999999		DW	9999H, 9999H, 9999H, 9999H, 9999H, 9999H, 9999H, 9999H	
00899	0426 9999999999		DW	9999H, 9999H, 9999H, 9999H, 9999H, 9999H, 9999H, 9999H	
00900	0436 9999999999		DW	9999H, 9999H, 9999H, 9999H, 9999H, 9999H, 9999H, 9999H	
00901	0446 00	DB#Save\$E:	DB	0	;E register
00902	0447 00	DB#Save\$I:	DB	0	;I register
00903	0448 00	DB#Save\$C:	DB	0	;C register
00904	0449 00	DB#Save\$B:	DB	0	;B register
00905	044A 00	DB#Save\$F:	DB	0	;Flags
00906	044B 00	DB#Save\$A:	DB	0	;A register
00907		DB#Stack:			;Debug stack area
00908					; The registers in the stack area are PUSHed
00909					; onto the stack and accessed directly.

Figure 10-2. (Continued)


```

00910 ;
00911 ; Register caption messages
00912 ;
00913 ; The table below, indexed by the Display%Code is used to access
00914 ; the register caption string.
00915 ;
00916 DB$Register$Captions:
00917 044C 7204 DW DB$F$RC ;Flags
00918 044E 7804 DW DB$A$RC ;A register
00919 0450 7A04 DW DB$B$RC ;B
00920 0452 7C04 DW DB$C$RC ;C
00921 0454 7E04 DW DB$D$RC ;D
00922 0456 8004 DW DB$E$RC ;E
00923 0458 8204 DW DB$H$RC ;H
00924 045A 8404 DW DB$L$RC ;L
00925 045C 8604 DW DB$BC$RC ;BC
00926 045E 8904 DW DB$DE$RC ;DE
00927 0460 8C04 DW DB$HL$RC ;HL
00928 0462 8F04 DW DB$SP$RC ;Stack pointer
00929 0464 9204 DW DB$M$RC ;Memory
00930 0466 A604 DW DB$B$BC$RC ;(BC)
00931 0468 AB04 DW DB$B$DE$RC ;(DE)
00932 046A B004 DW DB$B$HL$RC ;(HL)
00933 046C B504 DW DB$W$BC$RC ;(BC+1),(BC)
00934 046E C104 DW DB$W$DE$RC ;(DE+1),(DE)
00935 0470 CD04 DW DB$W$HL$RC ;(HL+1),(HL)
00936 ;
00937 0472 466C616773DB$F$RC: DB 'Flags',0 ;Flags
00938 0478 4100 DB$A$RC: DB 'A',0 ;A register
00939 047A 4200 DB$B$RC: DB 'B',0 ;B
00940 047C 4300 DB$C$RC: DB 'C',0 ;C
00941 047E 4400 DB$D$RC: DB 'D',0 ;D
00942 0480 4500 DB$E$RC: DB 'E',0 ;E
00943 0482 4800 DB$H$RC: DB 'H',0 ;H
00944 0484 4C00 DB$L$RC: DB 'L',0 ;L
00945 0486 424300 DB$BC$RC: DB 'BC',0 ;BC
00946 0489 444500 DB$DE$RC: DB 'DE',0 ;DE
00947 048C 484C00 DB$HL$RC: DB 'HL',0 ;HL
00948 048F 535000 DB$SP$RC: DB 'SP',0 ;Stack pointer
00949 0492 5374617274DB$M$RC: DB 'Start, End Address',0 ;Memory
00950 04A6 2842432900DB$B$BC$RC: DB '(BC)',0 ;(BC)
00951 04AE 2844452900DB$B$DE$RC: DB '(DE)',0 ;(DE)
00952 04B0 28484C2900DB$B$HL$RC: DB '(HL)',0 ;(HL)
00953 04B5 2842432B31DB$W$BC$RC: DB '(BC+1),(BC)',0 ;(BC+1),(BC)
00954 04C1 2844452B31DB$W$DE$RC: DB '(DE+1),(DE)',0 ;(DE+1),(DE)
00955 04CD 28484C2B31DB$W$HL$RC: DB '(HL+1),(HL)',0 ;(HL+1),(HL)
00956 ;
00957 ; Flags message
00958 ;
00959 04D9 43785A784DDB$Flags$Msg: DB 'CxZxMxExIx',0 ;Compatible with DDT's display
00960 ;
00961 ; Flags masks used to test user's flag byte
00962 ;
00963 DB$Flag$Masks:
00964 04E4 01 DB 0000$0001B ;Carry
00965 04E5 40 DB 0100$0000B ;Zero
00966 04E6 80 DB 1000$0000B ;Minus
00967 04E7 04 DB 0000$0100B ;Even parity
00968 04E8 10 DB 0001$0000B ;Interdigit carry (aux carry)
00969 04E9 00 DB 0 ;Terminator
01100 ;#
01101 ; DB$Init
01102 ; This routine initializes the debug package.
01103 ;
01104 ; DB$Init:
01105 IF DB$BIOS$IO ;Use BIOS for CONIN/CONOUT
01106 LHL 1 ;Get warm boot address from base
01107 ; page. H = BIOS jump vector page
01108 MVI L,09H ;Get CONIN offset in jump vector
01109 SHLD BIOS$CONIN + 1 ;Set up address
01110 MVI L,0CH ;Get CONOUT offset in jump vector
01111 SHLD BIOS$CONOUT + 1
01112 ENDIF
01113 ;
01114 ;Set up JMP instructions to receive control
01115 ; when an RST instruction is executed
01116 04EA 3EC3 MVI A,JMP ;Set JMP instructions at RST points

```

Figure 10-2. (Continued)

01117	04EC 322000	STA	RST4	
01118	04EF 322800	STA	RST5	
01119	04F2 211A08	LXI	H,DB#Input	;Address of fake input routine
01120	04F5 222100	SHLD	RST4 + 1	
01121	04F8 216C08	LXI	H,DB#Output	;Address of fake output routine
01122	04FB 222900	SHLD	RST5 + 1	
01123				
01124	04FE C9	RET		
01200		;	;	;
01201		;	DB#CONINU	
01202		;		This routine returns the next character from the console,
01203		;		but converting "a" to "z" to uppercase letters.
01204		;		;
01205		;	DB#CONINU:	
01206	04FF CD0505	CALL	DB#CONIN	;Get character from keyboard
01207	0502 C31B09	JMP	DB#A\$To\$Upper	;Fold to upper and return
01300		;	;	;
01301		;	DB#CONIN	
01302		;		This routine returns the next character from the console.
01303		;		According to the setting of equates, it uses simple
01304		;		polled I/O, the BDOS (function 2) or the BIOS.
01305		;		;
01306		;	Exit parameters	
01307		;		;
01308		;		A = character from console
01309		;		;
01310		;	DB#CONIN:	
01311		IF	DB#Polled\$I/O	;Simple polled input
01312		IN	DB#Status\$Port	;Check if incoming data
01313		ANI	DB#Input\$Ready	
01314		JZ	DB#CONIN	;No
01315		IN	DB#Data\$Port	;Input data character
01316		PUSH	PSW	;Save data character
01317		MOV	C,A	;Ready for output
01318		CALL	DB#CONOUT	;Echo it back
01319		POP	PSW	;Recover data character
01320		RET		
01321		ENDIF		
01322				
01323		IF	DB#BDOS\$I/O	;Use BDOS for input
01324	0505 0E01	MVI	C,DB#CONIN	;Read console
01325	0507 C30500	JMP	BDOS	;BDOS returns to our caller
01326		ENDIF		
01327				
01328		IF	DB#BIOS\$I/O	;Use BIOS for input
01329		JMP	BIOS#CONIN	;This was set up during BIOS
01330				; initialization
01331		ENDIF		
01332				
01400		;	;	;
01401		;	DB#CONOUT	
01402		;		This routine outputs the character in the C register to the
01403		;		console, using simple polled I/O, the BDOS or the BIOS.
01404		;		;
01405		;	Entry parameters	
01406		;		A = byte to be output
01407		;		;
01408		;	DB#CONOUT:	
01409	050A 3A0604	LDA	DB#Flag	;Check if debug output enabled
01410	050D B7	ORA	A	
01411	050E C8	RZ		;Ignore output if disabled
01412				
01413		IF	DB#Polled\$I/O	;Use simple polled output
01414		IN	DB#Status\$Port	;Check if ready for output
01415		ANI	DB#Output\$Ready	
01416		JZ	DB#CONOUT	;No
01417		MOV	A,C	;Get data byte
01418		OUT	DB#Data\$Port	
01419		RET		
01420		ENDIF		
01421				
01422		IF	DB#BDOS\$I/O	;Use BDOS for output
01423	050F 59	MOV	E,C	;Move into correct register
01424	0510 0E02	MVI	C,DB#CONOUT	
01425	0512 C30500	JMP	BDOS	;BDOS returns to our caller
01426		ENDIF		
01427				
01428		IF	DB#BIOS\$I/O	;Use BIOS for output

Figure 10-2. (Continued)

```

01429                MOV     A,C                ;Move into correct register
01430                JMP     BIOS*CONOUT        ;Set up during debug initialization
01431                ENDIF
01500                ;#
01501                ;
01502                ;     DB*On
01503                ;     This routine enables all debug output by setting the
01504                ;     DB*Flag nonzero.
01505                ;
01506                ;     DB*On:
01507                0515 F5        PUSH     PSW                ;Preserve registers
01508                0516 3EFF      MVI     A,OFFH
01509                0518 320604   STA     DB*Flag        ;Set control flag on
01510                051B F1        POP     PSW
01511                051C C9        RET
01600                ;#
01601                ;
01602                ;     DB*Off
01603                ;     This routine disables all debug output by setting the
01604                ;     DB*Flag to zero.
01605                ;
01606                ;     DB*Off:
01607                051D F5        PUSH     PSW                ;Preserve registers
01608                051E AF        XRA     A
01609                051F 320604   STA     DB*Flag        ;Clear control flag
01610                0522 F1        POP     PSW
01611                0523 C9        RET
01700                ;#
01701                ;
01702                ;     DB*Set*Pass
01703                ;     This routine sets the pass counter. Subsequent calls to DB*Pass
01704                ;     decrement the count, and when it reaches 0, debug output
01705                ;     is enabled.
01706                ;
01707                ;     Calling sequence
01708                ;
01709                ;             CALL     DB*Set*Pass
01710                ;             DW      Pass*Count*Value
01711                ;
01712                ;     DB*Set*Pass:
01713                0524 220904   SHLD   DB*Save*HL        ;Preserve user's HL
01714                0527 E1        POP     H                ;Recover return address
01715                0528 D5        PUSH   D                ;Preserve user's DE
01716                0529 5E        MOV     E,M
01717                052A 23        INX    H                ;Get LS byte of count
01718                052B 56        MOV     D,M
01719                052C 23        INX    H                ;Update pointer
01720                052D EB        XCHG   H,D                ;Get MS byte
01721                052E 220704   SHLD   DB*Pass*Count    ;HL points to return address
01722                0531 EB        XCHG   H,D                ;HL = pass counter
01723                0532 D1        POP     D                ;Set debug pass counter
01724                0533 E3        XTHL                      ;HL points to return address
01725                0534 C9        RET        ;Recover user's DE
01726                0534 C9        RET        ;Recover user's HL and set
01800                ;#                ; return address on top of stack
01801                ;
01802                ;     DB*Pass
01803                ;     This routine decrements the debug pass counter -
01804                ;     if the result is negative, it takes no further action.
01805                ;     If the result is zero, it sets the debug control flag nonzero
01806                ;     to enable debug output.
01807                ;
01808                ;     DB*Pass:
01809                0535 F5        PUSH   PSW                ;Save user's registers
01810                0536 E5        PUSH   H
01811                0537 2A0704   LHLD   DB*Pass*Count    ;Get pass count
01812                053A 2B        DCX    H
01813                053B 7C        MOV     A,H                ;Check if count now negative
01814                053C B7        ORA    A
01815                053D FA4705   JM     DB*Pass*x        ;Yes, take no further action
01816                0540 220704   SHLD   DB*Pass*Count    ;Save downated count
01817                0543 B5        ORA    L                ;Check if count now zero
01818                0544 CA4A05   JZ     DB*Pass*ED        ;Yes, enable debug
01819                ;
01820                ;     DB*Pass*x:
01821                0547 E1        POP     H                ;Recover user's registers
01822                0548 F1        POP     PSW
01822                0549 C9        RET

```

Figure 10-2. (Continued)

```

01823                ;
01824                DB*Pass*Ed:                ;Enable debug
01825    054A 3EFF      MVI        A,OFFH
01826    054C 320604   STA        DB*Flag      ;Set debug control flag
01827    054F C34705   JMP        DB*Pass*x
01900                ;#
01901                ;
01902                ;       DB*Display
01903                ;       This is the primary debug display routine.
01904                ;
01905                ;       Calling sequence
01906                ;
01907                ;       CALL    DB*Display
01908                ;       DB      Display*Code
01909                ;       DB      'Caption String',0
01910                ;
01911                ;       Display code identifies which register(s) are to be
01912                ;       displayed.
01913                ;
01914                ;       When the display code specifies a block of memory
01915                ;       the sequence is:
01916                ;
01917                ;       CALL    DB*Display
01918                ;       DB      Display*Code
01919                ;       DW      Start*Address,End*Address
01920                ;       DB      'Caption String',0
01921                ;
01922                ; DB*Display:
01923                ;
01924                ; DB*Display*Enabled:
01925    0552 220904   SHLD     DB*Save*HL      ;Save user's HL
01926                ;
01927    0555 E3       XTHL     ;Get return address from stack
01928    0556 220D04   SHLD     DB*Save*RA      ;This gets updated by debug code
01929    0559 E5       PUSH    H                ;Save return address temporarily
01930    055A 2B      DCX     H                ;Subtract 3 to address call instruction
01931    055B 2B      DCX     H                ; itself
01932    055C 2B      DCX     H
01933    055D 220F04   SHLD     DB*Call*Address ;Save actual address of CALL
01934    0560 E1       POP     H                ;Recover return address
01935                ;
01936    0561 F5       PUSH    PSW             ;Temporarily save flags to avoid
01937                ;       them being changed by DAD SP
01938    0562 210000   LXI     H,0            ;Preserve stack pointer
01939    0565 39      DAD     SP
01940    0566 23      INX     H                ;Correct for extra PUSH PSW needed
01941    0567 23      INX     H                ; to save the flags
01942    0568 220B04   SHLD     DB*Save*SP
01943    056B F1       POP     PSW             ;Recover flags
01944                ;
01945    056C 314C04   LXI     SP,DB*Stack     ;Switch to local stack
01946                ;
01947    056F F5       PUSH    PSW             ;Save other user's registers
01948    0570 C5      PUSH    B                ;The stack area is specially laid
01949    0571 D5      PUSH    D                ; out to access these registers
01950                ;
01951    0572 2A0D04   LHLD     DB*Save*RA      ;Get return address
01952    0575 7E      MOV     A,M            ;Get display code
01953    0576 321504   STA     DB*Display*Code
01954    0579 23      INX     H                ;Update return address
01955                ;
01956    057A FE18     CPI     DB*M            ;Check if memory to be displayed
01957    057C C29105   JNZ     DB*Not*Memory
01958    057F 5E      MOV     E,M            ;Get DE = start address
01959    0580 23      INX     H
01960    0581 56      MOV     D,M
01961    0582 23      INX     H
01962    0583 EB      XCHG    ;HL = start address
01963    0584 221104   SHLD     DB*Start*Address
01964    0587 EB      XCHG    ;HL -> end address
01965    0588 5E      MOV     E,M            ;Get DE = end address
01966    0589 23      INX     H
01967    058A 56      MOV     D,M
01968    058B 23      INX     H
01969    058C EB      XCHG    ;HL = end address, DE -> caption
01970    058D 221304   SHLD     DB*End*Address
01971    0590 EB      XCHG    ;HL -> caption string

```

Figure 10-2. (Continued)

```

01972          DB*Not*Memorys:
01973          ;
01974          ;       Output preamble and caption string
01975          ;       The format for everything except memory display is:
01976          ;
01977          ;       nnnn : Caption String : RC = vvvv
01978          ;       ^           ^           ^
01979          ;       Call Address           Value
01980          ;                               Register Caption (A, B, C,...)
01981          ;
01982          ;       A carriage return, line feed is output at the start of the
01983          ;       message - but NOT at the end.
01984          ;
01985          ;       Memory displays look like :
01986          ;
01987          ;       nnnn : Caption String : Start, End ssss, eeee
01988          ;       ssss : hh hh hh hh hh hh hh hh hh hh hh hh hh hh hh hh : cccc cccc cccc cccc
01989          ;
01990
01991          0591 E5          PUSH      H          ;Save pointer to caption string
01992          0592 CDC107     CALL       DB*CRLF        ;Display carriage return, line feed
01993          0595 CD7C07     CALL       DB*Display$CALLA ;Display DB*Call$Address in hex.
01994
01995          0598 E1          POP        H          ;Recover pointer to caption string
01996          ;
01997          0599 7E          DB*Display$Caption: ;HL -> caption string
01998          059A 23          MOV       A,M          ;Get character
01999          059B B7          INX       H          ;
02000          059C CAA805     ORA      A          ;Check if end of string
02001          ;
02002          059F E5          JZ       DB*End$Caption ;Yes
02003          ;
02004          05A0 4F          PUSH     H          ;Save string pointer
02005          05A1 CD0A05     MOV     C,A          ;Ready for output
02006          05A4 E1          CALL    DB*CONOUT     ;Display character
02007          05A5 C39905     POP     H          ;Recover string pointer
02008          ;
02009          ;                               JMP     DB*Display$Caption ;Go back for next character
02010          ;
02011          ; DB*End$Caption:
02012          05AB 220D04     SHLD    DB$Save$RA   ;Save updated return address
02013          ;
02014          05AB CDC807     CALL    DB$Colon     ;Display ' : '
02015          ;
02016          05AE 3A1504     LDA     DB*Display$Code ;Display register caption
02017          05B1 5F          MOV     E,A          ;Get user's display code
02018          05B2 1600     MVI    D,0          ;Make display code into word
02019          05B4 D5          PUSH   D          ;Save word value for later
02020          ;
02021          05B5 FE18     CPI    DB*$M         ;Memory display is a special case
02022          05B7 FACF05     JZ     DB*Display$Mem$Caption ;Yes
02023          ;
02024          05BA 214C04     LXI    H,DB*Register$Captions ;Make pointer to address in table
02025          05BD 19          DAD    D          ;HL -> word containing address of
02026          ;                               ; register caption
02027          05BE 5E          MOV     E,M          ;Get LS byte of address
02028          05BF 23          INX    H          ;
02029          05C0 56          MOV     D,M          ;DE -> register caption string
02030          05C1 EB          XCHG   ;HL -> register caption string
02031          05C2 CDEE07     CALL    DB*$MSG      ;Display message addressed by HL
02032          ;
02033          05C5 CDD607     CALL    DB*$MSGI     ;Display in-line message
02034          05C8 203D2000    DB     ', ',0
02035          05CC C3ED05     JMP     DB*$Select$Routine ;Go to correct processor
02036          ;
02037          ; DB*Display$Mem$Caption:
02038          ;                               ; The memory display requires a special
02039          ;                               ; caption with the start and end
02040          ;                               ; addresses
02041          ;                               ; Display ' : '
02042          05CF 219204     LXI    H,DB*$M$RRC ;
02043          05D2 CDEE07     CALL    DB*$MSG      ;Display specific caption
02044          05D5 CDC807     CALL    DB$Colon     ;Display ' : '
02045          ;
02046          05D8 2A1104     LHLD   DB*$Start$Address ;Display start address
02047          05DB CD8707     CALL    DB*$DHLH     ;Display HL in hex.
02048          ;
02049          05DE CDD607     CALL    DB*$MSGI     ;Display in-line message
02050          05E1 2C2000    DB     ', ',0
02051          ;
02052          05E4 2A1304     LHLD   DB*$End$Address ;Get end address

```

Figure 10-2. (Continued)

```

02049 05E7 CD8707      CALL  DB%DHLH      ;Display HL in hex.
02050 05EA CDC107      CALL  DB%CRLF     ;Display carriage return, line feed
02051                                     ;Drop into select routine
02052                                     DB$Select$Routine:
02053 05ED D1             POP    D          ;Recover word value Display$Code
02054 05EE 210A06        LXI   H,DB$Display$Table
02055 05F1 19           DAD    D          ;HL -> address of code to process
02056                                     ; display requirements
02057 05F2 5E           MOV    E,M       ;Get LS byte of address
02058 05F3 23           INX   H          ;Update pointer
02059 05F4 56           MOV    D,M       ;Get MS byte of address
02060 05F5 EB           XCHG          ;HL -> code
02061
02062 05F6 11FB05        LXI   D,DB$Exit  ;Fake link on stack
02063 05F9 D5           PUSH  D          ;
02064 05FA E9           PCHL          ;"CALL" display processor
02065
02066                                     DB$Exit:
02067 05FB D1           POP    D          ;Return to the user
02068 05FC C1           POP    B          ;Recover user's registers saved
02069 05FD F1           POP    PSW       ; on local debug stack
02070 05FE 2A0B04        LHLD  DB$Save$SP ;Revert to user's stack
02071 0601 F9           SPHL          ;
02072 0602 2A0D04        LHLD  DB$Save$RA ;Get updated return address (bypasses
02073                                     ; in-line parameters)
02074 0605 E3           XTHL          ;Replace on top of user's stack
02075 0606 2A0904        LHLD  DB$Save$HL ;Get user's HL
02076 0609 C9           RET           ;Transfer to correct return address
02077
02078                                     DB$Display$Table:
02079 060A 3006          DW    DP$F      ;Flags
02080 060C 5406          DW    DP$A      ;A register
02081 060E 5A06          DW    DP$B      ;B
02082 0610 6006          DW    DP$C      ;C
02083 0612 6606          DW    DP$D      ;D
02084 0614 6C06          DW    DP$E      ;E
02085 0616 7206          DW    DP$H      ;H
02086 0618 7806          DW    DP$L      ;L
02087 061A 7E06          DW    DP$BC     ;BC
02088 061C 8406          DW    DP$DE     ;DE
02089 061E 8A06          DW    DP$HL     ;HL
02090 0620 9006          DW    DP$SP     ;Stack pointer
02091 0622 9606          DW    DP$M      ;Memory
02092 0624 9C06          DW    DP$B$BC   ;(BC)
02093 0626 A206          DW    DP$B$DE   ;(DE)
02094 0628 A806          DW    DP$B$HL   ;(HL)
02095 062A B006          DW    DP$W$BC   ;(BC+1),(BC)
02096 062C B806          DW    DP$W$DE   ;(DE+1),(DE)
02097 062E C006          DW    DP$W$HL   ;(HL+1),(HL)
02098
02099                                     ;#
02100                                     ;
02101                                     ; Debug display processing routines
02102                                     ;
02103 DP$F:                                     ;Flags
02104                                     ;The flags are displayed in the same way that
02105                                     ; DDT uses: C1Z0M0E0I0
02106 0630 3A4A04        LDA   DB$Save$F ;Get flags
02107 0633 47           MOV    B,A       ;Preserve copy
02108 0634 21DA04        LXI   H,DB$Flags$Msg + 1 ;HL -> first 0/1 in message
02109 0637 11E404        LXI   D,DB$Flag$Masks ;DE -> table of flag mask values
02110
02111 DB$F$Next:
02112 063A 1A           LDAX  D          ;Get next flag mask
02113 063B B7           ORA   A          ;Check if end of table
02114 063C CA4E06        JZ    DB$F$Display ;Yes, display the results
02115
02116 063F A0           ANA   B          ;Check if this flag is set
02117 0640 3E31        MVI   A,'1'     ;Assume yes
02118 0642 C24706        JNZ   DB$F$NZ   ;Yes, it is set
02119 0645 3E30        MVI   A,'0'     ;No, it is clear
02120
02121 DB$F$NZ:
02122 0647 77           MOV    M,A       ;Store '0' or '1' in message text
02123 0648 23           INX   H          ;Update pointer to next 0/1
02124 0649 23           INX   H          ;
02125 064A 13           INX   D          ;Update flag mask pointer
02126 064B C33A06        JMP   DB$F$Next ;
02127 DB$F$Display:
02128 064E 21D904        LXI   H,DB$Flags$Msg ;Display results

```

Figure 10-2. (Continued)

```

02227 0651 C3EE07      JMP      DB$MESSG      ;Display message and return
02228
02229      ; DP#A:      ;A register
02230      LDA      DB$Save$A      ;Get saved value
02231 0654 3A4B04      JMP      DB$DAH        ;Display it and return
02232
02233      ; DP#B:      ;B
02234 065A 3A4904      LDA      DB$Save$B      ;Get saved value
02235 065D C39107      JMP      DB$DAH        ;Display it and return
02236
02237      ; DP#C:      ;C
02238 0660 3A4804      LDA      DB$Save$C      ;Get saved value
02239 0663 C39107      JMP      DB$DAH        ;Display it and return
02240
02241      ; DP#D:      ;D
02242 0666 3A4704      LDA      DB$Save$D      ;Get saved value
02243 0669 C39107      JMP      DB$DAH        ;Display it and return
02244
02245      ; DP#E:      ;E
02246 066C 3A4604      LDA      DB$Save$E      ;Get saved value
02247 066F C39107      JMP      DB$DAH        ;Display it and return
02248
02249      ; DP#H:      ;H
02250 0672 3A0A04      LDA      DB$Save$H      ;Get saved value
02251 0675 C39107      JMP      DB$DAH        ;Display it and return
02252
02253      ; DP#L:      ;L
02254 0678 3A0904      LDA      DB$Save$L      ;Get saved value
02255 067B C39107      JMP      DB$DAH        ;Display it and return
02256
02257      ; DP#BC:      ;BC
02258 067E 2A4804      LHL D  DB$Save$C      ;Get saved word value
02259 0681 C38707      JMP      DB$DHLH       ;Display it and return
02260
02261      ; DP#DE:      ;DE
02262 0684 2A4604      LHL D  DB$Save$E      ;Get saved word value
02263 0687 C38707      JMP      DB$DHLH       ;Display it and return
02264
02265      ; DP#HL:      ;HL
02266 068A 2A0904      LHL D  DB$Save$HL     ;Get saved word value
02267 068D C38707      JMP      DB$DHLH       ;Display it and return
02268
02269      ; DP#SP:      ;Stack Pointer
02270 0690 2A0B04      LHL D  DB$Save$SP     ;Get saved word value
02271 0693 C38707      JMP      DB$DHLH       ;Display it and return
02272
02273      ; DP#M:      ;Memory
02274 0696 2A1304      LHL D  DB$End$Address ;Increment end address to make
02275 0699 23          INX      H              ; arithmetic easier
02276 069A 221304      SHLD   DB$End$Address
02277
02278 069D 2A1104      LHL D  DB$Start$Address
02279 06A0 CD3A07      CALL   DB$M$Check$End ;Compare HL to End$Address
02280 06A3 DAD106      JC     DB$M$Address$OK ;End > start
02281 06A6 CDD607      CALL   DB$MESSGI      ;Error start > end
02282 06A9 OD0A2A20    DB      ODH,OAH,'** ERROR - Start Address > End **',0
02283 06CD C9          RET
02284
02285      ; DB$M$Next$Line:
02286 06CE CDC107      CALL   DB$CRLF        ;Output carriage return, line feed
02287
02288      ; DB$M$Address$OK:
02289 06D1 CDD607      CALL   DB$MESSGI      ;Bypass CR,LF for first line
02290 06D4 202000      DB      ' ',0          ;Indent line
02291 06D7 2A1104      LHL D  DB$Start$Address ;Get start of line address
02292 06DA CD8707      CALL   DB$DHLH        ;Display in hex
02293
02294 06DD CDC807      CALL   DB$Colon        ;Display ':'
02295
02296 06E0 2A1104      LHL D  DB$Start$Address
02297      ; DB$M$Next$Hex$Byte:
02298 06E3 E5          PUSH   H              ;Save memory address
02299 06E4 CDD007      CALL   DB$Blank        ;Output a blank
02300 06E7 E1          POP    H              ;Recover current byte address
02301 06E8 7E          MOV    A,M            ;Get byte from memory
02302 06E9 23          INX    H              ;Update memory pointer
02303 06EA E5          PUSH   H              ;Save for later
02304 06EB CD9107      CALL   DB$DAH        ;Display in hex.
02305 06EE E1          POP    H              ;Recover memory updated address

```

Figure 10-2. (Continued)

```

02305 06FF CD3A07 CALL DB#M#Check#End ;Compare HL vs.end address
02306 06F2 CAFE06 JZ DB#M#Display#ASCII ;Yes, end of area
02307 06F5 7D MOV A,L ;Check if at start of new line,
02308 06F6 E60F ANI 00001111B ; (is address XXX0H?)
02309 06F8 CAFE06 JZ DB#M#Display#ASCII ;Yes
02310 06FB C3E306 JMP DB#M#Next#Hex#Byte ;No, loop back for another
02311
02312 ;
02313 DB#M#Display#ASCII: ;Display bytes in ASCII
02313 06FE CDC807 CALL DB#Colon ;Display ':'
02314 0701 2A1104 LHLDB DB#Start#Address ;Start ASCII as beginning of line
02315 DB#M#Next#ASCII#Byte:
02316 0704 7E MOV A,M ;Get byte from memory
02317 0705 E5 PUSH H ;Save memory address
02318 0706 E67F ANI 01111111B ;Remove parity
02319 0708 4F MOV C,A ;Prepare for output
02320 0709 FE20 CPI ; ;Check if non-graphic
02321 070B D21007 JNC DB#M#Display#Char ;Char >= space
02322 070E 0E2E MVI C,'.' ;Display non-graphic as '.'
02323 DB#M#Display#Char:
02324 0710 FE7F CPI 7FH ;Check if DEL (may be non-graphic)
02325 0712 C21707 JNZ DB#M#Not#DEL ;No, it is graphic
02326 0715 0E2E MVI C,'.' ;Force to '.'
02327
02328 ;
02329 DB#M#Not#DEL:
02329 0717 CD0A05 CALL DB#CONOUT ;Display character
02330 071A E1 POP H ;Recover memory address
02331 071B 23 INX H ;Update memory pointer
02332 071C 221104 SHLD DB#Start#Address ;Update memory copy
02333 071F CD3A07 CALL DB#M#Check#End ;Check if end of memory dump
02334 0722 CA3707 JZ DB#M#Exit ;Yes, done
02335 0725 7D MOV A,L ;Check if end of line
02336 0726 E60F ANI 00001111B ; by checking address = XXX0H
02337 0728 CACE06 JZ DB#M#Next#Line ;Yes, start next line
02338 072B 7D MOV A,L ;Check if extra blank needed
02339 072C E603 ANI 00000011B ; if address is multiple of 4
02340 072E C20407 JNZ DB#M#Next#ASCII#Byte ;No -- go back for next character
02341 0731 CDD007 CALL DB#Blank ;Yes, output blank
02342 0734 C30407 JMP DB#M#Next#ASCII#Byte ;Go back for next character
02343
02344 ;
02345 DB#M#Exit:
02346 0737 C3C107 JMP DB#CRLF ;Output carriage return, line feed
02347 ; and return
02348
02349 ;
02350 DB#M#Check#End: ;Compares HL vs End#Address
02350 073A D5 PUSH D ;Save DE (defensive programming)
02351 073B EB XCHG ;DE = current address
02352 073C 2A1304 LHLDB DB#End#Address ;Get end address
02353 073F 7A MOV A,D ;Compare MS bytes
02354 0740 BC CMP H ;
02355 0741 C24607 JNZ DB#M#Check#End#X ;Exit now as they are unequal
02356 0744 7B MOV A,E ;Compare LS bytes
02357 0745 BD CMP L ;
02358 DB#M#Check#End#X:
02359 0746 EB XCHG ;HL = current address
02360 0747 D1 POP D ;Recover DE
02361 0748 C9 RET ;Return with condition flags set
02362
02363 ;
02364 DP#B#BC: ;(BC)
02364 0749 2A4804 LHLDB DB#Save#C ;Get saved word value
02365 074C 7E MOV A,M ;Get byte addressed by it
02366 074D C39107 JMP DB#DAH ;Display it and return
02367
02368 ;
02369 DP#B#DE: ;(DE)
02369 0750 2A4604 LHLDB DB#Save#E ;Get saved word value
02370 0753 7E MOV A,M ;Get byte addressed by it
02371 0754 C39107 JMP DB#DAH ;Display it and return
02372
02373 ;
02374 DP#B#HL: ;(HL)
02374 0757 2A0904 LHLDB DB#Save#HL ;Get saved word value
02375 075A 7E MOV A,M ;Get byte addressed by it
02376 075B C39107 JMP DB#DAH ;Display it and return
02377
02378 ;
02379 DP#W#BC: ;(BC+1),(BC)
02379 075E 2A4804 LHLDB DB#Save#C ;Get saved word value
02380 0761 5E MOV E,M ;Get word addressed by it
02381 0762 23 INX H ;

```

Figure 10-2. (Continued)


```

02382 0763 56          MOV     D,M
02383 0764 EB          XCHG          ;HL = word to be displayed
02384 0765 C38707     JMP     DB#DHLH ;Display it and return
02385
02386 ;
02386 DP#W#DE:          ;(DE+1),(DE)
02387 0768 2A4604     LHL     DB#Save#E ;Get saved word value
02388 0768 5E          MOV     E,M ;Get word addressed by it
02389 076C 23          INX     H
02390 076D 56          MOV     D,M
02391 076E EB          XCHG          ;HL = word to be displayed
02392 076F C38707     JMP     DB#DHLH ;Display it and return
02393
02394 ;
02394 DP#W#HL:          ;(HL+1),(HL)
02395 0772 2A0904     LHL     DB#Save#HL ;Get saved word value
02396 0775 5E          MOV     E,M ;Get word addressed by it
02397 0776 23          INX     H
02398 0777 56          MOV     D,M
02399 0778 EB          XCHG          ;HL = word to be displayed
02400 0779 C38707     JMP     DB#DHLH ;Display it and return
02401
02401 ;
02500 ;#
02501 ;
02501 DB#Display#CALLA
02502 ; This routine displays the DB#Call#Address in hexadecimal,
02503 ; followed by " : ".
02504 ;
02504 DB#Display#CALLA:
02506 077C E5          PUSH    H ;Save caller's HL
02507 077D 2A0F04     LHL     DB#Call#Address ;Get the call address
02508 0780 CD8707     CALL    DB#DHLH ;Display HL in hex.
02509 0783 E1          POP     H ;Recover caller's HL
02510 0784 C3C807     JMP     DB#Colon ;Display " : " and return
02511
02511 ;
02600 ;#
02601 ;
02602 ;
02602 DB#DHLH
02603 ; Display HL in hex.
02604 ;
02605 ;
02605 Entry parameters
02606 ;
02607 ; HL = value to be displayed
02608 ;
02609 DB#DHLH:
02610 0787 E5          PUSH    H ;Save input value
02611 0788 7C          MOV     A,H ;Get MS byte first
02612 0789 CD9107     CALL    DB#DAH ;Display A in hex.
02613 078C E1          POP     H ;Recover input value
02614 078D 7D          MOV     A,L ;Get LS byte
02615 078E C39107     JMP     DB#DAH ;Display it and return
02616
02616 ;
02700 ;#
02701 ;
02702 ;
02702 DB#DAH
02703 ; Display A register in hexadecimal
02704 ;
02705 ;
02705 Entry parameters
02706 ;
02707 ; A = value to be converted and output
02708 ;
02709 DB#DAH:
02710 0791 F5          PUSH    PSW ;Take a copy of the value to be converted
02711 0792 0F          RRC          ;Shift A right four places
02712 0793 0F          RRC
02713 0794 0F          RRC
02714 0795 0F          RRC
02715 0796 CDB407     CALL    DB#Nibble#To#Hex ;Convert LS 4 bits to ASCII
02716 0799 CDA0A5     CALL    DB#CONOUT ;Display the character
02717 079C F1          POP     PSW ;Get original value again
02718 079D CDB407     CALL    DB#Nibble#To#Hex ;Convert LS 4 bits to ASCII
02719 07A0 C3A0A5     JMP     DB#CONOUT ;Display and return to caller
02800
02800 ;#
02801 ;
02802 ;
02802 DB#CAH
02803 ; Convert A register to hexadecimal ASCII and store in
02804 ; specified address.
02805 ;
02806 ;
02806 Entry parameters
02807 ;

```

Figure 10-2. (Continued)

```

02808             ;           A = value to be converted and output
02809             ;           HL -> buffer area to receive two characters of output
02810             ;
02811             ;           Exit parameters
02812             ;
02813             ;           HL -> byte following last hex.byte output
02814             ;
02815             ; DB#CAH:
02816             07A3 F5      PUSH    PSW           ;Take a copy of the value to be converted
02817             07A4 0F      RRC          ;Shift A right four places
02818             07A5 0F      RRC
02819             07A6 0F      RRC
02820             07A7 0F      RRC
02821             07A8 CDB407  CALL    DB#Nibble$To$Hex      ;Convert to ASCII hex.
02822             07AB 77      MOV     M,A           ;Save in memory
02823             07AC 23      INX     H           ;Update pointer
02824             07AD F1      POP     PSW          ;Get original value again
02825             07AE CDB407  CALL    DB#Nibble$To$Hex      ;Convert to ASCII hex.
02826             07B1 77      MOV     M,A           ;Save in memory
02827             07B2 23      INX     H           ;Update pointer
02828             07B3 C9      RET
02900             ;#
02901             ;
02902             ;           Minor subroutines
02903             ;
02904             ;
02905             ;           DB#Nibble$To$Hex
02906             ;           This is a minor subroutine that converts the least
02907             ;           significant four bits of the A register into an ASCII
02908             ;           hex. character in A and C
02909             ;
02910             ;           Entry parameters
02911             ;
02912             ;           A = nibble to be converted in LS 4 bits
02913             ;
02914             ;           Exit parameters
02915             ;
02916             ;           A,C = ASCII hex. character
02917             ;
02918             ; DB#Nibble$To$Hex:
02919             07B4 E60F      ANI    00001111B        ;Isolate LS four bits
02920             07B6 C630      ADI    '0'           ;Convert to ASCII
02921             07B8 FE3A      CPI    '9' + 1       ;Compare to maximum
02922             07BA DABF07    JC     DB#NTH$Numeric ;No need to convert to A -> F
02923             07BD C607      ADI    7             ;Convert to a letter
02924             ; DB#NTH$Numeric:
02925             07BF 4F      MOV    C,A           ;For convenience of other routines
02926             07C0 C9      RET
02927             ;
02928             ;
02929             ;           DB#CRLF
02930             ;           Simple routine to display carriage return, line feed.
02931             ;
02932             ; DB#CRLF:
02933             ;
02934             07C1 CDD607    CALL    DB#MSGI       ;Display in-line message
02935             07C4 OD0A00    DB     ODH,0AH,0
02936             07C7 C9      RET
02937             ;
02938             ;           DB#Colon
02939             ;           Simple routine to display ':'.
02940             ;
02941             ; DB#Colon:
02942             07C8 CDD607    CALL    DB#MSGI       ;Display in-line message
02943             07CB 203A2000  DB     ':',0
02944             07CF C9      RET
02945             ;
02946             ;           DB#Blank
02947             ;           Simple routine to display ''.
02948             ;
02949             ; DB#Blank:
02950             07D0 CDD607    CALL    DB#MSGI       ;Display in-line message
02951             07D3 2000      DB     '',0
02952             07D5 C9      RET
03100             ;#
03101             ;
03102             ;           Message processing subroutines

```

Figure 10-2. (Continued)

```

03103      ;
03104      ;      DB$MESSG (message in-line)
03105      ;      Output null-byte terminated message that follows the
03106      ;      CALL to MSGOUTI
03107      ;
03108      ;      Calling sequence
03109      ;
03110      ;      CALL      DB$MESSG
03111      ;      DB      'Message',0
03112      ;      ... next instruction
03113      ;
03114      ;      Exit parameters
03115      ;      HL -> instruction following message
03116      ;
03117      ;
03118      DB$MESSG:
03119      ;
03120      ;      ;Get return address of stack, save
03121      ;      ; user's HL on top of stack
03122      ;      ;HL -> message
03123      07D6 E3      XTHL
03124      07D7 F5      PUSH    PSW      ;Save all user's registers
03125      07D8 C5      PUSH    B
03126      07D9 D5      PUSH    D
03127      DB$MESSG$Next:
03128      07DA 7E      MOV     A,M      ;Get next data byte
03129      07DB 23      INX     H      ;Update message pointer
03130      07DC B7      ORA     A      ;Check if null byte
03131      07DD C2E507   JNZ     DB$MESSGIC ;No, continue
03132      07E0 D1      POP     D      ;Recover user's registers
03133      07E1 C1      POP     B
03134      07E2 F1      POP     PSW
03135      07E3 E3      XTHL
03136      ;      ;Recover user's HL from stack, replacing
03137      ;      ; it with updated return address
03138      ;      ;Return to address after 00-byte
03139      ;      ; after in-line message
03140      DB$MESSGIC:
03141      07E5 E5      PUSH    H      ;Save message pointer
03142      07E6 4F      MOV     C,A      ;Ready for output
03143      07E7 CD0A05   CALL    DB$CONOUT
03144      07EA E1      POP     H      ;Recover message pointer
03145      07EB C3DA07   JMP     DB$MESSG$Next ;Go back for next char.
03146      ;
03147      ;      DB$MESSG
03148      ;      Output null-byte terminated message
03149      ;
03150      ;      Calling sequence
03151      ;
03152      ;      MESSAGE:      DB      'Message',0
03153      ;
03154      ;      LXI     H,MESSAGE
03155      ;      CALL    DB$MESSG
03156      ;
03157      ;      Exit parameters
03158      ;      HL -> null byte terminator
03159      ;
03160      ;
03161      DB$MESSG:
03162      07EE F5      PUSH    PSW      ;Save user's registers
03163      07EF C5      PUSH    B
03164      07F0 D5      PUSH    D
03165      DB$MESSG$Next:
03166      07F1 7E      MOV     A,M      ;Get next byte for output
03167      07F2 B7      ORA     A      ;Check if 00-byte terminator
03168      07F3 CA0008   JZ     DB$MESSG$X ;Exit
03169      07F6 23      INX     H      ;Update message pointer
03170      07F7 E5      PUSH    H      ;Save updated pointer
03171      07F8 4F      MOV     C,A      ;Ready for output
03172      07F9 CD0A05   CALL    DB$CONOUT
03173      07FC E1      POP     H      ;Recover message pointer
03174      07FD C3F107   JMP     DB$MESSG$Next ;Go back for next character
03175      ;
03176      DB$MESSG$X:
03177      0800 D1      POP     D      ;Recover user's registers
03178      0801 C1      POP     B
03179      0802 F1      POP     PSW

```

Figure 10-2. (Continued)

```

03180 0803 C9          RET
03300                ;#
03301                ;
03302                ;   Debug input routine
03303                ;
03304                ;   This routine helps debug code in which input instructions
03305                ;   would normally occur. The opcode of the IN instruction
03306                ;   must be replaced by a value of 0E7H (RST 4).
03307                ;
03308                ;   This routine picks up the port number contained in the byte
03309                ;   following the RST 4, converts it to hexadecimal, and
03310                ;   displays the message:
03311                ;
03312                ;       Input from port XX :
03313                ;
03314                ;   It then accepts two characters (in hex.) from the keyboard,
03315                ;   converts these to binary in A, and then returns control
03316                ;   to the byte following the port number
03317                ;
03318                ;   *****
03319                ;   WARNING - This routine uses both DB#CONOUT and BDOS calls
03320                ;   *****
03321                ;
03322 0804 496E707574DBIN#Message: DB      'Input from Port '
03323 0814 5858203A20DBIN#Port:   DB      'XX : ',0
03324                ;
03325                ;
03326                ;   DB#Input:
03327 081A 220904        SHLD   DB#Save#HL      ;Save user's HL
03328 081D E1           POP    H                ;Recover address of port number
03329 081E 2B          DCX    H                ;Backup to point to RST
03330 081F 220F04      SHLD   DB#Call#Address ;Save for later display
03331 0822 23          INX    H                ;Restore to point to port number
03332                ;Note: A need not be preserved
03333 0823 7E          MOV    A,M            ;Get port number
03334 0824 23          INX    H                ;Update return address to bypass port number
03335 0825 220D04      SHLD   DB#Save#RA      ;Save return address
03336 0828 C5          PUSH   B                ;Save remaining registers
03337 0829 D5          PUSH   D                ;
03338 082A F5          PUSH   PSW           ;Save port number for later
03339                ;
03340                ;
03341 082B CDB108        CALL   DB#Flag#Save#On ;Save current state of debug flag
03342                ; and enable debug output
03343                ;
03344 082E CDC107        CALL   DB#CRLF          ;Display carriage return, line feed
03345 0831 CD7C07        CALL   DB#Display#CALLA;Display call address
03346 0834 F1          POP    PSW           ;Recover port number
03347 0835 211408      LXI    H,DBIN#Port
03348 0838 CDA307      CALL   DB#CAH          ;Convert to hex. and store in message
03349 083B 210408      LXI    H,DBIN#Message ;Output prompting message
03350 083E CDEE07      CALL   DB#MSG
03351 0841 0E02        MVI    C,2            ;Get 2 digit hex. value
03352 0843 CDCF08      CALL   DB#GHV          ;Returns value in HL
03353 0846 7D          MOV    A,L            ;Get just single byte
03354                ;
03355 0847 CDBF08        CALL   DB#Flag#Restore ;Restore debug output to previous state
03356                ;
03357 084A D1          POP    D                ;Recover registers
03358 084B C1          POP    B                ;
03359 084C 2A0904      LHLD  DB#Save#HL      ;Get previous HL
03360 084F E5          PUSH   H                ;Put on top of stack
03361 0850 2A0D04      LHLD  DB#Save#RA      ;Get return address
03362 0853 E3          XTHL
03363 0854 C9          RET
03500                ;#
03501                ;
03502                ;   Debug output routine
03503                ;
03504                ;   This routine helps debug code in which output instructions
03505                ;   would normally occur. The opcode of the OUT instruction
03506                ;   must be replaced by a value of 0EFH (RST 5).
03507                ;
03508                ;   This routine picks up the port number contained in the byte
03509                ;   following the RST 5, converts it to hexadecimal, and
03510                ;   displays the message:
03511                ;

```

Figure 10-2. (Continued)

```

03512                ;           Output to port XX : AA
03513                ;
03514                ;           where AA is the contents of the A register prior to the
03515                ;           RST 5 being executed.
03516                ;           Control is then returned to the byte following the port number.
03517                ;
03518                ;           *****
03519                ;           WARNING - This routine uses both DB*CONOUT and BDOS calls
03520                ;           *****
03521                ;
03522                ;
03523 0855 4F75747075DB0$Message:  DB      'Output to Port '
03524 0864 5858203A20DB0$Port:    DB      'XX : '
03525 0869 414100 DB0$Value:      DB      'AA',0
03526                ;
03527                ;
03528                ; DB$Output:
03529 086C 220904                SHLD   DB$Save$HL      ;Save user's HL
03530 086F E1                    POP    H              ;Recover address of port number
03531 0870 2B                    DCX   H              ;Backup to point to RST
03532 0871 220F04                SHLD   DB$Call$Address ;Save for later display
03533 0874 23                    INX   H              ;Restore to point at port number
03534 0875 324B04                STA   DB$Save$A       ;Preserve value to be output
03535 0878 7E                    MOV   A,M            ;Get port number
03536 0879 23                    INX   H              ;Update return address to bypass port number
03537 087A 220D04                SHLD   DB$Save$RA     ;Save return address
03538 087D C5                    PUSH  B              ;Save remaining registers
03539 087E D5                    PUSH  D              ;
03540 087F F5                    PUSH  PSW            ;Save port number for later
03541                ;
03542 0880 CDB108                CALL  DB$Flag$Save$On ;Save current state of debug flag
03543                ;           ; and enable debug output
03544                ;
03545 0883 CDC107                CALL  DB$CRLF         ;Display carriage return, line feed
03546 0886 CD7C07                CALL  DB$Display$CALLA ;Display call address
03547 0889 F1                    POP    PSW           ;Recover port number
03548 088A 216408                LXI   H,DB0$Port
03549 088D CDA307                CALL  DB$CAH         ;Convert to hex.and store in message
03550                ;
03551 0890 3A4B04                LDA   DB$Save$A
03552 0893 216908                LXI   H,DB0$Value
03553 0896 CDA307                CALL  DB$CAH         ;Convert to hex.and store in message
03554                ;
03555 0899 215508                LXI   H,DB0$Message ;Output prompting message
03556 089C CDEE07                CALL  DB$MSG
03557                ;
03558 089F CDBF08                CALL  DB$Flag$Restore ;Restore debug flag to previous state
03559                ;
03560 08A2 D1                    POP    D              ;Recover registers
03561 08A3 C1                    POP    B              ;
03562 08A4 2A0904                LHLD  DB$Save$HL     ;Get previous HL
03563 08A7 E5                    PUSH  H              ;Put on top of stack
03564 08A8 2A0D04                LHLD  DB$Save$RA     ;Get return address
03565 08AB E3                    XTHL ;TOS = return address, HL = previous value
03566 08AC 3A4B04                LDA   DB$Save$A
03567 08AF C9                    RET
03700                ;#
03701                ;
03702                ;           DB$Flag$Save$On
03703                ;           This routine is only used for DB$IN/OUT.
03704                ;           It saves the current state of the debug control flag,
03705                ;           D$Flag, and then enables it to make sure that
03706                ;           DB$IN/OUT output always goes out.
03707                ;
03708 08B0 00                DB$Flag$Previous:  DB      0      ;Previous flag value
03709                ;
03710                ; DB$Flag$Save$On:
03711 08B1 F5                    PUSH  PSW            ;Save caller's registers
03712 08B2 3A0604                LDA   DB$Flag       ;Get current value
03713 08B5 32B008                STA   DB$Flag$Previous ;Save it
03714 08B8 3EFF                MVI   A,0FFH        ;Set flag
03715 08BA 320604                STA   DB$Flag
03716 08BD F1                    POP    PSW
03717 08BE C9                    RET
03800                ;#
03801                ;

```

Figure 10-2. (Continued)

```

03802                ;         DB$Flag$Restore
03803                ;         This routine is only used for DB$IN/OUT.
03804                ;         It restores the debug control flag, DB$Flag, to
03805                ;         its former state.
03806                ;
03807                DB$Flag$Restore:
03808                08BF F5                PUSH    PSW
03809                08C0 3AB008           LDA     DB$Flag$Previous      ;Get previous setting
03810                08C3 320604           STA     DB$Flag              ;Set debug control flag
03811                08C6 F1                POP     PSW
03812                08C7 C9                RET
03813
03814                ;
03900                ;#
03901                ;
03902                ;         Get hex. value
03903                ;
03904                ;         This subroutine outputs a prompting message, and then reads
03905                ;         the keyboard in order to get a hexadecimal value.
03906                ;         It is somewhat simplistic in that the first non-hex value
03907                ;         terminates the input. The maximum number of digits to be
03908                ;         converted is specified as an input parameter. If more than the
03909                ;         maximum number is entered, only the last four are significant.
03910                ;
03911                ;*****
03912                ;                     W A R N I N G
03913                ;         DB$GHV will always use the BDOS to perform a read console
03914                ;         function (#10). Be careful if you use this routine from
03915                ;         within an executing BIOS.
03916                ;*****
03917                ;
03918                ;         Entry parameters
03919                ;
03920                ;         HL -> 00-byte terminated message to be output
03921                ;         C = number of hexadecimal digits to be input
03922                ;
03923                ;
03924                DB$GHV$Buffer:         ;Input buffer for console characters
03925                DB$GHV$Max$Count:     ;
03926                08C8 00                DB     0                        ;Set to the maximum number of chars.
03927                ;                     ; to be input
03928                DB$GHV$Input$Count:   ;
03929                08C9 00                DB     0                        ;Set by the BDOS to the actual number
03930                ;                     ; of chars. entered
03931                DB$GHV$Data$Bytes:    ;
03932                08CA                DS     5                        ;Buffer space for the characters
03933                ;
03934                ;
03935                DB$GHV:
03936                08CF 79                MOV     A,C                      ;Get maximum characters to be input
03937                08D0 FE05             CPI     5                        ;Check against maximum count
03938                08D2 DAD708           JC     DB$GHV$Count$OK          ;Carry set if A < 5
03939                08D5 3E04             MVI    A,4                      ;Force to only four characters
03940                DB$GHV$Count$OK:
03941                08D7 32C808           STA     DB$GHV$Max$Count        ;Set up maximum count in input buffer
03942                08DA CDEE07           CALL   DB$MSG                   ;Output prompting message
03943                08DD 11C808           LXI    D,DB$GHV$Buffer         ;Accept characters from console
03944                08E0 0E0A             MVI    C,B$READCONS           ;Function code
03945                08E2 CD0500           CALL   BDOS
03946                ;
03947                08E5 0E02             MVI    C,B$CONOUT              ;Output a line feed
03948                08E7 1E0A             MVI    E,0AH
03949                08E9 CD0500           CALL   BDOS
03950                ;
03951                08EC 210000           LXI    H,0                      ;Initial value
03952                08EF 11CA08           LXI    D,DB$GHV$Data$Bytes     ;DE -> data characters
03953                08F2 3AC908           LDA     DB$GHV$Input$Count     ;Get count of characters input
03954                08F5 4F              MOV     C,A                      ;Keep count in C
03955                DB$GHV$Loop:
03956                08F6 0D              DCR    C                        ;Downdate count
03957                08F7 F8              RM
03958                08F8 1A              LDAX   D                        ;Return when all done (HL has value)
03959                08F9 13              INX   D                        ;Get next character from buffer
03960                08FA CD1B09           CALL   DB$A$To$Upper           ;Update buffer pointer
03961                08FD FE30             CPI    '0'                      ;Convert A to uppercase if need be
03962                08FF D8              RC                               ;Check if less than 0
03963                0900 FE3A             CPI    '9' + 1                 ;Yes, terminate
03964                0902 DA1009           JC     DB$GHV$Hex$Digit        ;Check if > 9
                                ;No, it must be numeric

```

Figure 10-2. (Continued)

```

03965 0905 FE41      CPI    'A'           ;Check if < 'A'
03966 0907 D8        RC          ;Yes, terminate
03967 0908 FE47      CPI    'F' + 1      ;Check if > 'F'
03968 090A D0        RNC          ;Yes, terminate
03969 090B D637      SUI    'A' - 10     ;Convert A through F to numeric
03970 090D C31209    JMP    DB$GHV$Shift$Left$4 ;Combine with current result
03971
03972                ;
03972                DB$GHV$Hex$Digit:
03973 0910 D630      SUI    '0'           ;Convert to binary
03974                DB$GHV$Shift$Left$4:
03975 0912 29         DAD    H             ;Shift HL left four bits
03976 0913 29         DAD    H
03977 0914 29         DAD    H
03978 0915 29         DAD    H
03979 0916 85         ADD    L             ;Add binary value in LS 4 bits of A
03980 0917 6F        MOV    L,A           ;Put back into HL total
03981 0918 C3F608    JMP    DB$GHV$Loop  ;Loop back for next character
04100
04101                ;#
04102                ;
04103                ;   A to upper
04104                ;   Converts the contents of the A register to an uppercase
04105                ;   letter if it is currently a lowercase letter
04106                ;
04107                ;   Entry parameters
04108                ;
04109                ;       A = character to be converted
04110                ;
04111                ;   Exit parameters
04112                ;
04113                ;       A = converted character
04114                ;
04114                DB$A$To$Upper:
04115 091B FE61      CPI    'a'           ;Compare to lower limit
04116 091D D8        RC          ;No need to convert
04117 091E FE7B      CPI    'z' + 1      ;Compare to upper limit
04118 0920 D0        RNC          ;No need to convert
04119 0921 E65F      ANI    5FH          ;Convert to uppercase
04120 0923 C9        RET

```

Figure 10-2. Debug subroutines (continued)

```

B>ddt fig10-2.hex<cr>
DDT VERS 2.0
NEXT PC
0924 0000
-g100<cr>

0116 : Flags : Flags = C1Z0M1E110
0120 : A Register : A = AA
012F : B Register : B = BB
013E : C Register : C = CC
014D : D Register : D = DD
015C : E Register : E = EE
016B : H Register : H = FF
017A : L Register : L = 11
0189 : Memory Dump #1 : Start, End Address : 0108, 0128
0108 : 05 3E AA 01 CC BB 11 EE : .>*. L;.n
0110 : DD 21 11 FF B7 37 CD 52 05 00 46 6C 61 67 73 00 : !!.. 77MR ..F1 ags.
0120 : CD 52 05 02 41 20 52 65 67 : MR.. A Re g
01A0 : Memory Dump #2 : Start, End Address : 0100, 011F
0100 : 31 6B 03 CD EA 04 CD 15 05 3E AA 01 CC BB 11 EE : 1k.M j.M. .>*. L;.n
0110 : DD 21 11 FF B7 37 CD 52 05 00 46 6C 61 67 73 00 : !!.. 77MR ..F1 ags.
01B7 : Memory Dump #3 : Start, End Address : 0101, 0100
** ERROR - Start Address > End **
01CE : Memory Dump #4 : Start, End Address : 0100, 0100
0100 : 31 : 1

```

Figure 10-3. Console output from debug testbed run

```

01E5 : BC Register : BC = BBCC
01F5 : DE Register : DE = DDEE
0205 : HL Register : HL = FF11
0215 : SP Register : SP = 0369
022E : Byte at (BC) : (BC) = BC
023F : Byte at (DE) : (DE) = DE
0250 : Byte at (HL) : (HL) = F1
026A : Word at (BC) : (BC+1), (BC) = 0B0C
027B : Word at (DE) : (DE+1), (DE) = 0D0E
028C : Word at (HL) : (HL+1), (HL) = 0F01
Debug output has been re-enabled.
This message should display 5 times
This message should display 5 times
This message should display 5 times
This message should display 5 times
This message should display 5 times
032B : Input from Port 11 : aa

032D : Output to Port 22 : AA

```

Figure 10-3. Console output from debug tested run (continued)

containing all of the symbols in your program, along with their respective addresses. Once the program has been loaded by SID, you can refer to the memory image of your program not by address, but by the actual symbol name from your source code. SID also supports the “pass count” concept when using breakpoints.

ZSID (Z80 Symbolic Debug)

This is the Z80 CPU's version of SID. The mini-assembler/disassembler uses Zilog instruction mnemonics rather than those used by Intel.

Bringing Up CP/M for the First Time

It is much harder to bring up CP/M on a new computer system than to debug an enhanced version on a system already running CP/M. You will often find yourself staring at a programmatic “brick wall” with no adequate debugging tools to assist you.

For example, you install the CP/M system on a diskette (using another CP/M-based computer system), put the diskette into the new computer, and press the RESET button. The disk head loads on the disk, and then — nothing! You cannot use any programs such as DDT or SID because you do not yet have CP/M up and running on the new computer. Or can you?

The answer is, wherever possible, debug the code for the new machine on an existing CP/M system. You may have to “fake” some aspects of the new bootstrap or BIOS so that the act of testing it on the host machine does not interact with the CP/M already running on it.

This scheme permits you to be fairly sure of your program logic before loading the diskette into the new machine. It will help pin down problems caused by hardware problems on the new computer.

The hardest situation of all is if you have only the new computer and the release diskettes from Digital Research. Your only option is to find a way of reading the CP/M image on the release diskette into memory, hand patch in new console and disk drivers (not a trivial task), write the patched image back onto a diskette, and resort to Orville Wright testing.

If you value your time, it is always more cost-effective to use another system with CP/M already installed. This is true even if the two systems do not have the same diskette format. You can still do the bootstrap and build the CP/M image on the host machine. Then download the image directly into the memory of the new machine and write it out to a diskette.

This *downloading* process does require, however, that the new computer have a read-only memory (ROM) monitor program. Depending on the capability of this ROM monitor program, you may have to hand patch into the new machine's memory a primitive "download" program that reads 8-bit characters from a serial port, stacking them up in memory and returning control to the monitor program when you press a keyboard character on the new machine's console. In fact, some ROM monitor programs have a downloading program built in.

Debugging the CP/M Bootstrap Loader

The CP/M bootstrap loader, as you may recall, is written on one of the outermost tracks on a diskette or hard disk. On a standard 8-inch single-sided, single-density diskette, CP/M's bootstrap loader is stored on the first sector of the first track. The loader is brought into memory by firmware that gets control of the CPU when you turn your machine on or press the RESET button.

The bootstrap has to be compact, as the diskette space on which it is stored is limited: no more than 128 bytes for standard 8-inch diskettes. This tends to rule out the use of the debug subroutines already described, so you have to fall back to more primitive techniques.

Testing the Bootstrap Under CP/M

A bootstrap is best developed on a CP/M-based system. The task is easiest of all if you already have CP/M running on your new machine and are simply preparing an enhanced version of the bootstrap loader. In this case, you can test most of the code as though it were a user program running in the transient program area (TPA).

Most bootstraps get loaded into memory at location 0000H, so at the front of the code to be debugged you must put a temporary origin line that reads

```
ORG      100H
```

If you omit this and ask DDT to load the HEX file output by the assembler, it will load at the true origin, 0000H, and wipe out the contents of the base page for the version of CP/M that you are running. This will cause a system crash; you will have to press the RESET button and reload CP/M. When this happens, DDT does not tell you directly that anything is amiss; it just displays a “?” after your request to load the HEX file. You will discover that the system has “gone away” only when you try to do something else.

You also will need to adjust the addresses into which the bootstrap tries to load the CP/M image. If you do not, you will overwrite the version of CP/M presently running.

With these adjustments made, you can load the bootstrap under DDT and watch it execute, confirming that it does load the correct image into the correct addresses for debugging and transfer control to the BIOS jump vector. When everything appears to be functioning correctly, use the IF instruction to disable the debug code, reassemble the bootstrap, and write it onto a diskette. Then put the diskette into drive A and press RESET.

Was the Bootstrap Loaded?

At this point you must establish whether the bootstrap is being loaded into memory when the machine is turned on or RESET is pressed. The best way of doing this, and one that you can leave in place permanently, is to output a sign-on message as soon as the loader gets control. This requires hardware set up to prepare the USART (Universal Synchronous/Asynchronous Receive/Transmit) chip to output data, although some manufacturers write this initialization code into the firmware that loads the bootstrap. A suitable sign-on message would be the following:

```
CP/M Bootstrap Loader : Vn 1.0 11/18/82
```

If you do not see this message, assume that control is *not* being transferred to the bootstrap loader. This will be useful in the future if someone should call you with a complaint that CP/M cannot be loaded. If this message does not appear, they probably do not have CP/M on the disk.

Did the Bootstrap Load CP/M?

This is a harder question to answer than whether the bootstrap itself has been loaded, especially if the bootstrap loader sign-on is displayed and then the system crashes. A sign-on message early in the BIOS cold boot processing can confirm the correct transfer of control into the BIOS.

If the problems with the bootstrap program are severe, you may have to adapt the memory-dump debugging subroutine, dumping the contents of memory to the console in order to see what information the bootstrap loader is placing in memory. Display 100H bytes starting from the front of the BIOS jump vector. This

table has an immediately recognizable pattern of 0C3H values every three bytes.

You should also check to see that the bootstrap is loading the correct number of sectors from the disk into memory. If it loads too few, CP/M may sign on only to crash a few moments later because it attempts either to execute code or access a constant at the end of the BIOS. If the bootstrap loads too many sectors from the disk, the excess may “wrap around” the top of memory and overwrite the bootstrap itself, down at location 0000H, before it has completed its task. In this case, you would see only the sign-on for the bootstrap, not for the BIOS.

Debugging the BIOS

Rather than try to debug the BIOS as a single piece of code, debug it as a series of separate functional modules.

Notwithstanding current “top-down” philosophies of dealing with overall structure first, it can be quicker to debug the low-level subroutines in a device driver first. This gives you a solid base on which to build.

The BIOS can be divided up into its constituent modules as follows:

Character input

- Interrupt service
- Non-interrupt service

Character output

Interrupt routines

- Real time clock
- Watchdog timers

Disk drivers

- High-level (deblocking)
- Low-level (physical I/O)

Plan to write a *testbed* program for each of these modules. This testbed code serves two purposes; first, it provides a means of transferring control into the module under test in a controlled way. Second, it includes the necessary modules or dummy modules to “fool” the module under test into responding as if it were running in a complete BIOS under CP/M.

Using the testbed, you can check every part of the module’s logic except the part that may be time-critical. Problems caused by timing, such as interrupts disabled for too long or code that is too slow or too fast for a particular peripheral controller chip, tend to show up only when you are testing on the final hardware and when you are running your new BIOS under CP/M.

What You Should Test for in the BIOS

Describing fully how to debug each module in the BIOS could fill several books. Remember that you are trying to establish the *absence* of errors using a technique that, by its very nature, tends to show only their *presence*.

There are two basic approaches to debugging. One is the plodding method, checking every aspect of the code to ensure that every feature really does work. The second is to try to do something useful with the code.

Plan to use both. Start with the plodding method, testing each feature under control of the testbed until you are sure that it is working *in vitro*. When all of the BIOS modules have been tested individually, build a CP/M system and try to do some useful work with it. Trying to use the system for actual work testing *in vitro* can be a good test.

Feature Checklist

Make a list of the specific features included in the various BIOS modules. Then devise specific test sequences that will show that each of the features is working correctly.

The same testbed code can often test all of the features of a driver module. If it cannot, create a new testbed for the more exotic features.

Keep the testbed routines. Experience shows that they are most often needed shortly after you have erased them. Even after you have tested the BIOS, the testbed routines will come in handy if you decide to enhance a particular driver later on. You can extract the driver code from the BIOS, glue it together with the testbed, and test the new feature code in isolation from the BIOS.

The following sections show example testbeds for the various drivers, along with example checklists. These checklists were used to test the example BIOS routines shown in earlier chapters.

Character Drivers

Figure 10-4 shows the code for an example testbed routine for character I/O drivers in the BIOS. This code would be followed by the actual character I/O drivers, exactly as they would appear in the BIOS except that all IN and OUT instructions would be replaced with RST 4's and 5's respectively (see Figure 10-2) so that you could enter input values and inspect output values on the console.

This example contains the initialization code for the debug package shown in Figure 10-2 and the code setting up an RST 6 used to "fake" incoming character interrupts.

The main testbed loop consists of a faked incoming character interrupt followed by optional calls to CONIN or CONOUT, the return of control to DDT, or a loop back to fake another character interrupt. You can only return control to DDT if you used DDT to load the testbed and driver programs in the first place.

```

;
; Testbed for character I/O drivers in the BIOS
;
; The complete source file consists of three components:
;
; 1. The testbed code shown here
; 2. The character I/O drivers destined for the BIOS
; 3. The debug package shown in Figure 10-2.
;
FFFF = TRUE EQU OFFFHH
0000 = FALSE EQU NOT TRUE

FFFF = DEBUG EQU TRUE ;For conditional assembly of RST
; instructions in place of IN and
; OUT instructions in the drivers
0030 = RST6 EQU 30H ;Use RST 6 for fake incoming character
; interrupt
0100 ORG 100H

START:
0100 31D101 LXI SP,Test*Stack ;Use a local stack
0103 CDD101 CALL DB*Init ;Initialize the debug package
0106 3EC3 MVI A,JMP ;Set up RST 6 with JMP opcode
0108 323000 STA RST6
010B 21D101 LXI H,Character*Interrupt ;Set up RST 6 JMP address
010E 223100 SHLD RST6 + 1
;
; Make repeated entry to character interrupt routine
; to ensure that characters can be captured and stored in
; an input buffer
;
;
Testbed*Loop:
0111 3EAA MVI A,0AAH ;Set registers to known pattern
0113 01CCBB LXI B,0BCCCH
0116 11EEDD LXI D,0DDEEH
0119 2111FF LXI H,OFF11H
011C F7 RST 6 ;Fake interrupt for incoming character

011D CDD101 CALL DB*MSGI ;Display in-line message
0120 0D0A456E74 DB 0DH,0AH,'Enter I to Input Char., 0 to Output, D to enter
0152 444454203A DB 'DDT : ',0

0159 CDD101 CALL DB*CONINU ;Get uppercase character
015C FE49 CPI 'I' ;CONIN?
015E CA7201 JZ Go*CONIN
0161 FE44 CPI 'D' ;DDT?
0163 CA6E01 JZ Go*DDT
0166 FE4F CPI 'O' ;CONOUT?
0168 CA9101 JZ Go*CONOUT
016B C31101 JMP Testbed*Loop ;Loop back to interrupt again

Go*DDT:
016E FF RST 7 ;Enter DDT (RST 7 set up by DDT)
016F C31101 JMP Testbed*Loop

Go*CONIN:
0172 CDD101 CALL CONST ;Get console status
0175 CA1101 JZ Testbed*Loop ;No data waiting
0178 CDD101 CALL CONIN ;Get data from buffer

017B CDD101 CALL DB*Display ;Display character returned
017E 02 DB DB*A ; in A register
017F 434F4E494E DB 'CONIN returned',0

018E C37201 JMP Go*CONIN ;Repeat CONIN loop until no chars.
; waiting
;
Go*CONOUT:
0191 CDD101 CALL CONST ;Get console status
0194 CA1101 JZ Testbed*Loop ;No data waiting
0197 CDD101 CALL CONIN
019A 4F MOV C,A ;Ready for output
019B CDD101 CALL CONOUT ;Output to console
019E C39101 JMP Go*CONOUT ;Repeat while there is still data
;
01A1 9999999999 DW 9999H,9999H,9999H,9999H,9999H,9999H,9999H,9999H
01B1 9999999999 DW 9999H,9999H,9999H,9999H,9999H,9999H,9999H,9999H
01C1 9999999999 DW 9999H,9999H,9999H,9999H,9999H,9999H,9999H,9999H

```

Figure 10-4. Testbed for character I/O drivers in the BIOS

```

Test$Stack:
;
;       Dummy routines for those shown in other figures
;
;       BIOS routines (Figure 8-10)
;
CONST:           ;BIOS console status
CONIN:           ;BIOS console input
CONOUT:          ;BIOS console output;
Character$Interrupt: ;Interrupt service routine for incoming chars.
;
;       Debug routines (Figure 10-2)
;
DB$Init:         ;Debug initialization
DB$MSGI:         ;Display message in-line
DB$CONINU:       ;Get uppercase character from keyboard
DB$Display:      ;Main debug display routine
DB%A EQU        02 ;Display code for DB$Display

0002 =

```

Figure 10-4. Testbed for character I/O drivers in the BIOS (continued)

Executing an RST 7 without using DDT will cause a system crash, as DDT sets up the necessary JMP instruction at location 0038H in the base page.

The faked incoming character interrupt transfers control directly to the interrupt service routine in the BIOS (see the example in Figure 8-10, line 04902, label Character\$Interrupt). This reads the status ports of each of the character devices; you can enter the specific status byte values that you want. If you enter a value that indicates that a data character is “incoming,” you will be prompted for the actual 8-bit data value to be “input.” You can make the interrupt service routine appear to be inputting characters and stacking characters up in the input buffer. For debugging purposes, reduce the size of the input buffer to eight bytes. Making it larger means you will have to input more characters to test the buffer threshold logic. To check the interrupt service routine, you will pass through the main testbed loop doing nothing but faking incoming character interrupts and entering status and data values. The data characters will then be stacked up in the input buffer.

To check the correct functioning of the interrupt service routines, you can stay in control with DDT from the outset. Alternatively, you can just use DDT to load the testbed/driver HEX file, loop around inputting several characters, and then request that the testbed return control to DDT. Then you can use DDT to inspect the contents of the device table(s) and input buffers.

Another possibility is to create debugging routines that display the contents of the device table in a meaningful way, with each field captioned like this:

```

DEVICE TABLE 0
  Status Port      81   Data Port      80
  Output Ready    01   Input Ready    02
  DTR high        40
  Reset Int. Prt  D8   Reset Int. Val. 20
  :
  :
  Status Byte 1
    Output Suspended
    Output Xon Enabled
  :

```

```

:
Buffer Base 0E8C
Put Offset   05      Get Offset   01
Char. Count  04      Control Count  00
Data Buffer
41 42 43 44 45 00 00 00

```

This display device table routine will require a fair amount of effort to code and debug—but it will pay dividends. You can obtain a complete “snapshot” of the device table without having to decode hexadecimal memory dumps and individual bits. Constant values in the device tables are also displayed, so that if a bug in your code corrupts the table, you will know about it immediately.

The next section shows examples of the specific tests you need to make, along with a description of the strategy you can use.

Interrupt Service Routine Checklist In a functioning BIOS, control is transferred to the interrupt service module whenever an incoming character causes an interrupt. In the example BIOS in Figure 8-10 (line 4900), the code scans each character device in turn to determine which one is causing the interrupt.

When you are debugging the interrupt service routines using the “fake” input/output instructions, you will have to enter specific status byte values. Refer to the device table declarations in Figure 8-10, line 1500, to determine what values you must enter to make the service routine think that an incoming character is arriving or that data terminal ready (DTR) is high or low.

Start the debugging process using the first device table. Then repeat the tests on the other device tables.

The following is a checklist of features that should be checked in debugging the interrupt service routine:

Are all registers restored correctly on exit from the interrupt servicing?

Using DDI, start execution from the beginning of the testbed. Set a breakpoint (with the G100,nnnn command) to get control back immediately before the CALL Character\$Interrupt. Use the X command to display all of the registers, and then, by using the G,nnnn command, you set a breakpoint at the instruction that immediately follows the CALL Character\$Interrupt. The character drivers will prompt you for the status values. Enter 00 (which indicates that no character is incoming). Display the registers again—their values should be the same. Remember to check the value of the stack pointer and the amount of the stack area that has been used.

NOTE: Do not be too surprised if you lose control of the machine when you first try this test. You may have some fundamental logic errors initially. If the system crashes, reset it, reload CP/M, and then start the test again. This time, rather than setting the second breakpoint at the instruction following the CALL Character\$Interrupt, venture down into the Character\$Interrupt code and go through the code a few instructions

at a time, setting breakpoints before any instructions that could cause a transfer of control. Find out how far you are getting into the driver before it either jumps off into space or settles into a loop.

Does the service routine push a significant number of bytes onto the stack after an interrupt has occurred?

When you get control back after the CALL Character\$Interrupt, use the D (dump) command to dump the stack area's memory on the console. Check how far down the stack came by looking for the point where the constants that used to fill the stack area are overwritten by other data.

The example BIOS in Figure 8-10 saves only the contents of the HL register pair on the pre-interrupt stack. It then switches over to a private BIOS stack to save the contents of the rest of the registers and service the interrupt.

Are data characters added to the input buffer correctly?

"Input" a noncontrol character via the Character\$Interrupt routine. Then check the contents of the appropriate device table. The character count and the put offset should both be set to one. Then check the contents of the input buffer itself; does it contain the character that you "input?"

Are control characters added to the input buffer correctly?

"Input" a control character such as 01H. Do not use ETX, ACK, XON, or XOFF (03H, 06H, 11H, and 13H, respectively); these may cause side effects if you have errors in the protocol handling logic. Check that the character is stored in the next byte of the input buffer and that the character and control counts are set to two and one, respectively. The put offset should also be set to two.

When the input buffer full threshold is reached, does the driver output the correct protocol character?

Set the first status byte in the first device table to enable input XON or RTS protocol, or both. Then go round the main testbed loop putting characters into the input buffer. Check the console display to see if the drivers output the correct values when the buffer is almost full (the default threshold is when five bytes remain). The driver should then drop the RTS line or output an XOFF character or both, according to the input protocol that you enabled.

When the input buffer is completely full, does the driver respond correctly?

This is an extension of the test above. Input one more character than can fit into the buffer. Check to see that the drivers do not stack the character into the input buffer and that a BELL character (07H) is output to the data port.

Are protocol characters XON/XOFF recognized and the necessary control flags set or reset?

Reload the testbed and drivers. Set the status byte to enable the output XON/XOFF protocol. Then use the Character\$Interrupt routine to input an XOFF character (13H). Check to see that the XOFF character has not been put into the input buffer. Instead, the status byte should be set to indicate that output has indeed been suspended.

Input an XON and check to see that the output suspended flag has been reset.

Does the driver detect and reset hardware errors correctly?

Proceed as though you were going to input a character into the input buffer, but instead enter a status byte value that indicates that a hardware error has occurred (enter the value given in the device table for DT\$Detect\$Error\$Value).

Check that the driver detects the error status and outputs the correct error-reset value to the appropriate control port.

Non-interrupt Service Routine Checklist In a “live” BIOS, non-interrupt service routines are accessed via the CONIN and CONST entry points in the BIOS jump vector. During debugging, the testbed can call the CONIN and CONST code directly.

Is input redirection functioning? Does control arrive in the driver with the correct device table selected?

This is best tested directly with DDT. Use the Gnnnn,bbbb command to transfer control into the CONIN code with a breakpoint at the RET instruction at the end of the Select\$Device\$Table routine (see Figure 8-10, line 04400). Check that the DE register pair is pointing at device table 0. If it is not, you will have to restart the test. Use the Tn command to make DDT trace through the Select\$Device\$Table subroutine to find the bug.

Are characters returned correctly from the buffer?

Use the testbed to “input” a character or two. Then use the testbed to make several entries into CONIN. Check the characters returned from the buffer.

Are the data character and control character counts correctly decremented?

After each character has been removed from the buffer by CONIN, use DDT to examine the device table and check that the data character and control character counts have been decremented correctly. Also check that the get pointer has moved up the input buffer.

When the buffer “almost empty” threshold is reached, does the driver emit the correct protocol character or manipulate the request to send (RTS) line correctly?

Use DDT to enable the input RTS or XON protocol or both. Then input characters into the input buffer until it reaches the buffer full threshold (the

default is when only five spare bytes remain in the buffer). Confirm that “buffer almost full” processing occurs. Then make repetitive calls to CONIN to flush data out of the buffer. Check that the “buffer emptying” processing occurs when the correct threshold is reached. For RTS protocol, the driver should output a raise RTS value to the specified RTS control port. For XON, the driver should output an XON character to the data port (after first having read the status port to ensure that the hardware can output the character).

Does the driver handle buffer “wraparound” correctly?

Input characters to the input buffer until it becomes completely full. Then make a single CONIN call to remove the first character from the buffer. Follow this by inputting one more character to the buffer. Check that the get pointer is set to one and the put pointer set to zero.

Next, make successive CONIN calls to empty the buffer. Then input one more character to the buffer. Check that this last character is put into the first byte of the input buffer.

Can the driver handle “forced input” correctly?

Using DDT, set the forced input pointer to point to a 00-byte-terminated string; for example, use one of the function key decode default strings. (In Figure 8-10, the forced input pointer is initialized to point to a “startup string”—this is declared at the beginning of the configuration block at line 00400.)

Using DDT, call the CONST routine and check that it returns with A = 0FFH (indicating that there appears to be input data waiting).

Make successive calls to CONIN and confirm that the data bytes in the forced input string are returned. Check that the forcing of input ends when the 00H-byte is detected.

Does the console status routine operate correctly when it checks for data characters in the buffer, control characters in the buffer, and forced input?

Input a single noncontrol character, such as 41H, into the input buffer. Using DDT, check that the second status byte in the device table has the fake type-ahead flag set to zero. Call the CONST routine—it should return with A = 0FFH (meaning that there is data in the buffer). Then set the fake type-ahead bit in the second status byte and call CONST again. It should return with A = 00H (meaning that there is now “no data” in the buffer). Input a single control character into the buffer. Now CONST should return with A = 0FFH because there is a control character in the buffer.

Does the driver recognize escape sequences incoming from keyboard function keys?

This is a difficult feature to test when the real time clock routine is not running. The driver uses the watchdog timer to wait until all characters in

the escape sequence have arrived. You will therefore have to modify the code in CONIN so that the watchdog timer appears to time out immediately, rather than waiting for the real time clock to tick. To make this change, refer to Figure 8-10, line 2200; this is the start of the CONIN routine. Look for the label CONIN\$Wait\$For\$Delay. A few instructions later there is a JNZ CONIN\$Wait\$For\$Delay. Using DDT, set all three bytes of this JNZ to 00H.

Then, using the testbed, input the complete escape sequence into the input buffer. For example, input hexadecimal values 1B, 4F, 51 (ESCAPE, O, P), which correspond to the characters emitted on a VT-100 terminal when FUNCTION KEY 1 (PF1) is pressed.

Next, use the testbed to make successive calls to CONIN. You should see the text associated with the function key (FUNCTION KEY 1, LINE FEED) being returned by CONIN.

Repeat this test using different function key sequences, including a sequence that does not correspond to any of the preset function keys. Check that the escape sequence itself is returned by CONIN without being changed into another string.

Can the driver differentiate between a function key and the same escape sequence generated by discrete key strokes?

This is almost the same test as above. Make the same patch to the CONIN code, only this time do not enter the complete escape sequence into the buffer. Enter only the hex characters 1B and 4F. Make sure that the CONIN routine does not substitute another string in place of this quasi-escape sequence.

This test only mimics the results of manually entering an escape sequence. You could not press the keys on a terminal fast enough to get all three characters into the input buffer within the time allowed by the watchdog timer.

Character Output Checklist *Can the driver output a character?*

The CONOUT option in the testbed calls CONIN first to get a character. To start with, you may want to use DDT to set the C register to some graphic ASCII character such as 41H (A), and transfer control into CONOUT directly. Check that CONOUT reads the USART's status, waits for the output ready value, and then outputs the data to the data port. Note that the testbed will output all characters waiting in the input buffer (or forced input) when you select its CONOUT option. This is a convenience for advanced testing of the drivers—for initial testing you may want to modify the testbed to make only one call to CONIN and CONOUT and then return to the top of the testbed loop.

Does the driver suspend output when a protocol control flag indicates that output is to be suspended?

Using DDT, set the status byte in the device table to enable output XON/XOFF protocol. Then input an XOFF character and confirm that the output suspended bit in the status byte is set. Output a single character, and using DDT, confirm that the driver will remain in a status loop waiting for the output suspended bit to be cleared. Clear the bit using DDT and check that the character is output correctly.

When using ETX/ACK protocol, does the driver output an ETX after the specified number of characters have been output, then indicate that output is suspended?

For debugging purposes, alter the ETX message count value in the device table to three bytes. Then output three bytes of data via CONOUT. Check that the driver sends an ETX character (03H) after the three bytes have been output and that the output suspended flag in the status byte has been set.

Then input an ACK character (06H). Check that this character is not stored in the input buffer and that the output suspended flag is cleared.

Does the driver recognize and output escape sequences?

Input an ESCAPE, "t" (1BH, 74H) into the input buffer. Then output them via CONOUT. Using DDT, check that the CONOUT routine recognizes that an escape sequence is being output and selects the correct processing routine. In this case, the forced input pointer should be set to point at the ASCII time of day in the configuration block.

Does each of the escape sequence processors function correctly? Can the time and date be set to specified values using escape sequences?

Repeat the test above using all of the other escape sequences to make sure that they can be recognized and that they function correctly.

Real Time Clock Routines

A separate testbed program, shown in Figure 10-5, is used to check these routines. It calls the interrupt service routine directly to simulate a real time clock "tick," and then displays the time of day in ASCII on the console.

As you can see, the testbed makes a call into the debug package's initialization routine, DB\$Init, and then uses an RST 6 to generate fake clock "ticks."

There is a JMP instruction in the testbed that bypasses a call to Set\$Watchdog. Remove this JMP, either by editing it out or by using DDT to change it to NO OPERATIONS (NOP, 00H) when you are ready to test the watchdog routines.

Real Time Clock Test Checklist *Is the clock running at all?*

Using DDT, trace through the interrupt service routine logic. Check that the seconds are being updated.

```

;
; Testbed for real time clock driver in the BIOS.
;
; The complete source file consists of three components:
;
; 1. The testbed code shown here
; 2. The real time clock driver destined for the BIOS.
; 3. The debug package shown in Figure 10-2.
;
FFFF = TRUE EQU OFFFH
0000 = FALSE EQU NOT TRUE

FFFF = DEBUG EQU TRUE ;For conditional assembly of RST
; instructions in place of IN and
; OUT instructions in the drivers.
;Use RST 6 for fake clock tick.

0030 = RST6 EQU 30H

0100
START: ORG 100H

0100 318B01 LXI SP,Test$Stack ;Use local stack
0103 CD8B01 CALL DB$Init ;Initialize the debug package
0106 3EC3 MVI A,JMP ;Set up RST 6 with JMP opcode
0108 323000 STA RST6
010B 218B01 LXI H,RTC$Interrupt ;Set up RST 6 JMP address
010E 223100 SHLD RST6 + 1

0111 C31D01 JMP Testbed$Loop ;<=== REMOVE THIS JMP WHEN READY TO
; TEST WATCHDOG ROUTINES

0114 013200 LXI B,50 ;50 ticks before timeout
0117 214201 LXI H,WD$Timeout ;Address to transfer to
011A CD8B01 CALL Set$Watchdog ;Set the watchdog timer
;
;
; Make repeated entry to RTC interrupt routine
; to ensure that clock is correctly updated
;
Testbed$Loop:
011D 3EAA MVI A,0AAH ;Set registers to known pattern
011F 01CCBB LXI B,0BCCCH
0122 11EEDD LXI D,0DDEEH
0125 2111FF LXI H,OFF11H
0128 F7 RST 6 ;Fake interrupt clock

0129 CD8B01 CALL DB$MSGI ;Display in-line message
012C 436C6F636B DB 'Clock =',0

0134 218B01 LXI H,Time$In$ASCII ;Get address of clock in driver
0137 CD8B01 CALL DB$MSG ;Display current clock value
; (Note: Time$In$ASCII already has
; a line feed character in it)

013A CD8B01 CALL DB$MSGI ;Display in-line message
013D 0D00 DB ODH,0 ;Carriage return

013F C31D01 JMP Testbed$Loop
;
; Control arrives here when the watchdog timer times
; out
WD$Timeout:
0142 CD8B01 CALL DB$MSGI
0145 0D0A576174 DB ODH,0AH,'Watchdog timed out',0
015A C9 RET ;Return to watchdog routine
;
015B 9999999999 DW 9999H,9999H,9999H,9999H,9999H,9999H,9999H,9999H
016B 9999999999 DW 9999H,9999H,9999H,9999H,9999H,9999H,9999H,9999H
017B 9999999999 DW 9999H,9999H,9999H,9999H,9999H,9999H,9999H,9999H

Test$Stack:
;
; Dummy routines for those shown in other figures
;
; BIOS routines (Figure 8-10)
;
RTC$Interrupt: ;Interrupt service routine for clock tick
Set$Watchdog: ;Set watchdog timer
Time$In$ASCII: ;ASCII string of HH:MM:SS, LF, 0
;
; Debug routines (Figure 10-2)
;
DB$Init: ;Debug initialization
DB$MSGI: ;Display message in-line
DB$MSG: ;Display message

```

Figure 10-5. Testbed for real-time-clock driver in the BIOS

Are the hours, minutes, and seconds carrying over correctly?

Let the testbed code run at full speed. You should see the time being updated on the console display—although it will be updated much more rapidly than real time.

Use DDT to set the minutes to 58 and then let the clock run again. Does it correctly show the hour and reset the minutes to 00? Then set the hours to 11 and the minutes to 58 and let the clock run. Do minutes carry over into hours and are hours reset to 0?

Repeat these tests with the clock update constants set for 24-hour format.

Is the clock interrupt service routine restoring the registers correctly?

Using DDT, check that the registers are still set correctly on return from the clock interrupt service routine.

How much of a load on the pre-interrupt stack is the service routine imposing?

Check the “low water mark” of the preset values remaining in the testbed stack area to see how much of a load the interrupt service routine is imposing on the stack.

Can the watchdog timer be set to a nonzero value? Can it be set back to zero?

Using the second part of the testbed, call the Set\$Watchdog routine, and then monitor the testbed's execution as the watchdog timer times out. Check that the registers and stack pointer are set correctly when control is transferred to the timeout routine. Also check that control is returned properly from this routine, and thence from the interrupt service routine.

Disk Drivers

It is only feasible to check the low-level disk drivers in isolation from a real BIOS, as the BDOS interface to the deblocking code is very difficult to simulate. The testbed shown in Figure 10-6 serves only as a time-saver. It does not test the interface to the subroutines. Use DDT to set up the disk, track, and sector numbers, and then monitor the calls into SELDSK, SETTRK, SETSEC, SETDMA, and the read/write routines.

Unless you have the same disk controller on the host system as you do on the target machine, you will have to use the fake input/output system described earlier in this chapter, rather than attempt to read and write on real disks.

You can see that the testbed, after initializing the debugging package, makes calls to SELDSK, SETTRK, SETSEC, and SETDMA. It then calls a low-level read or write routine. The low-level routine called depends on which driver you wish to debug. For the standard floppy diskette driver shown in Figure 8-10, use Read\$No\$Deblock and Write\$No\$Deblock. For the 5 1/4-inch diskettes, use Read\$Physical and Write\$Physical. You will have to use DDT to set up some of the variables required by the low-level drivers that would normally be set up by the deblocking code.

```

;       Testbed for disk I/O drivers in the BIOS
;
;       The complete source file consists of three components:
;
;       1. The testbed code shown here
;       2. The Disk I/O drivers destined for the BIOS
;       3. The debug package shown in Figure 10-2.
;
FFFF =   TRUE   EQU   OFFFFH
0000 =   FALSE  EQU   NOT TRUE

FFFF =   DEBUG  EQU   TRUE           ;For conditional assembly of RST
;       ; instructions in place of IN and
;       ; OUT instructions in the drivers.

0100                ORG   100H
START:
0100 314704          LXI   SP,Test#Stack ;Use a local stack
0103 CD4704          CALL  DB#Init      ;Initialize the debug package
;
;       Make calls to SELDSK, SETTRK, SETSEC and SETDMA,
;       then either a read or write routine.
;
Testbed$Loop:
0106 314704          LXI   SP,Test#Stack ;Use local stack
;
0109 3A1202          LDA   Logical$Disk ;Set up for SELDSK call
010C 4F             MOV   C,A
010D CD4704          CALL  SELDSK
;
0110 CD4704          CALL  DB$Display   ;Display return value in HL
0113 14             DB   DB$HL
0114 53454C4453      DB   'SELDISK returned',0
;
0124 223201          SHLD  DPH$Start   ;Set up to display disk parameter header
0127 111000          LXI   D,16        ;Compute end address
012A 19             DAD   D
012B 223401          SHLD  DPH$End    ;Store into debug call
;
012E CD4704          CALL  DB$Display   ;Display DPH
0131 18             DB   DB#M        ;Memory
;
DPH$Start:
0132 0000           DW   0
DPH$End:
0134 0000           DW   0
0136 53656C6563     DB   'Selected DPH',0
;
0143 2A1302          LHLD  Track       ;Call SETTRK
0146 E5             PUSH H
0147 C1             POP  B             ;SETTRK needs track in BC
0148 CD4704          CALL  SETTRK
;
014B 3A1502          LDA   Sector     ;Call SETSEC
014E 4F             MOV   C,A
014F CD4704          CALL  SETSEC
;
0152 011702          LXI   B,Test#Buffer ;Set DMA address
0155 CD4704          CALL  SETDMA
0158 3A1602          LDA   Write$Disk  ;Check if reading or writing
015B B7             ORA   A
015C C2D101          JNZ   Test$Write
;
015F CD4704          CALL  Read$No$Deblock ;*** or Read$Physical depending on which
;*** drivers you are testing
0162 CD4704          CALL  DB$Display   ;Display return code
0165 02             DB   DB#A
0166 5465737420     DB   'Test Read returned',0
;
0179 CD0102          CALL  Check$Ripple ;Check if ripple pattern in buffer
017C CA0601          JZ    Testbed$Loop ;Yes, it is correct
;
017F CD4704          CALL  DB#MSGI     ;Indicate problem
0182 14             DB   DB$HL
0183 526970706C     DB   'Ripple pattern incorrect. HL -> failure.',0
;
01AC CD4704          CALL  DB$Display   ;Display test buffer:
01AF CD1800          CALL  DB#M        ;Memory
01B2 1702          DW   Test$Buffer

```

Figure 10-6. Testbed for disk I/O drivers in the BIOS

```

01B4 0002      DW      Test$Buffer$Size
01B6 436F6E7465  DB      'Contents of Test$Buffer',0

01CE C30601    JMP      Testbed$Loop

Test$Write:
01D1 CDF201    CALL    Fill$Ripple      ;Fill the test buffer with ripple pattern
01D4 CD4704    CALL    Write$No$Deblock;*** or Write$Physical depending on which
                    ;*** drivers you are testing

01D7 CD4704    CALL    DB$Display      ;Display return code
01DA 02        DB      DB$A
01DB 5465737420 DB      'Test Write returned',0

01EF C30601    JMP      Testbed$Loop

Fill$Ripple:
                    ;Fills the Test$Buffer with a pattern
                    ; formed by putting into each byte, the
                    ; least significant 8-bits of the byte's
                    ; address.

01F2 010002    LXI    B,Test$Buffer$Size
01F5 211702    LXI    H,Test$Buffer

FR$Loop:
01F8 75        MOV    M,L      ;Set pattern value into buffer
01F9 23        INX    H      ;Update buffer pointer
01FA 0B        DCX    B      ;Down date count
01FB 79        MOV    A,C      ;Check if count zero
01FC B0        ORA    B      ;
01FD C2F801    JNZ    FR$Loop ;Repeat until zero
0200 C9        RET

;
; Check$Ripple:
                    ;Check that the buffer is filled with the
                    ; correct ripple pattern
                    ; Returns with zero status if this is true,
                    ; nonzero status if the ripple is not
                    ; correct. HL point to the offending byte
                    ; (which should = L)

0201 010002    LXI    B,Test$Buffer$Size
0204 211702    LXI    H,Test$Buffer

CR$Loop:
0207 7D        MOV    A,L      ;Get correct value
0208 BE        CMP    M      ;Compare to that in the buffer
0209 C0        RNZ                    ;Mismatch, nonzero already indicated
020A 23        INX    H      ;Update buffer pointer
020B 0B        DCX    B      ;Downdate count
020C 79        MOV    A,C      ;Check count zero
020D B0        ORA    B      ;
020E C20702    JNZ    CR$Loop ;Repeat until zero
0211 C9        RET ;Zero flag will already be set

;
; Testbed variables
;
0212 00        Logical$Disk: DB      0      ;A = 0, B = 1,...
0213 0000     Track:      DW      0      ;Disk track number
0215 00        Sector:     DB      0      ;Disk sector number
0216 00        Write$Disk: DB      0      ;NZ to write to disk
;
0200 =        Test$Buffer$Size EQU    512 ;<=== Alter as required
0217          Test$Buffer: DS      Test$Buffer$Size
;
0417 9999999999 DW      9999H,9999H,9999H,9999H,9999H,9999H,9999H,9999H
0427 9999999999 DW      9999H,9999H,9999H,9999H,9999H,9999H,9999H,9999H
0437 9999999999 DW      9999H,9999H,9999H,9999H,9999H,9999H,9999H,9999H

Test$Stack:
;
; Dummy routines for those shown in other figures
;
; BIOS routines (Figure 8-10)
;
SELDSK:        ;Select logical disk
SETTRK:        ;Set track number
SETSEC:        ;Set sector number
SETDMA:        ;Set DMA address
Read$No$Deblock: ;Driver read routines
Read$Physical:
Write$No$Deblock: ;Driver write routines
Write$Physical:

```

Figure 10-6. (Continued)


```

;
;   Debug routines (Figure 10-2)
;
DB*Init:           ;Debug initialization
DB*MSGI:          ;Display message in-line
DB*Display:       ;Main debug display routine
0002 = DB*A      EQU    02   ;Display codes for DB*Display
0014 = DB*HL     EQU    20
0018 = DB*M      EQU    24

```

Figure 10-6. Testbed for disk I/O drivers in the BIOS (continued)

Before issuing the write call, the testbed fills the disk buffer with a known pattern. This pattern is checked on return from a read operation.

For both reading and writing, the testbed shows the contents of the A register. If you have added the enhanced disk error handling described in the previous chapter, the return value in A must *always* be zero.

Disk Driver Checklist *Does SELDSK return the correct address and set up the required system variables?*

Check that the correct disk parameter header address is returned for legitimate logical disks. Check, too, that it returns an address of 0000H for illegal disks.

Check that any custom processing, such as setting the disk type and deblocking requirements from extra bytes on the disk parameter blocks, is performed correctly.

Does the SETTRK and SETSEC processing function correctly?

Using DDT, check that the correct variables are set to the specified values.

Does the driver read in the spare-sector directory correctly?

Set up to execute a physical read and, using DDT, trace the logic of the READ entry point. Check that the spare-sector directory would be loaded into the correct buffer. If you are using fake input/output, use DDT to patch in a typical spare-sector directory with two or three “spared-out” sectors.

Does the driver produce the correct spare sector in place of a bad one?

Continuing with the physical read operation, check that, for “good” track/sectors, the sector-sparing logic returns the original track and sector number, and for “bad” track/sectors, it substitutes the correct spare track and sector. If you are using sector skipping, check that the correct number of sectors is skipped.

Can a sector be read in from the disk?

Continuing further with the physical read, check that the correct sector is read from the specified disk and track. If you are using real I/O (as

opposed to faking it), the “ripple pattern” set by the testbed can be used, or you can fill the disk buffer area with some known pattern (using DDT's F command) so you can tell if any data gets read in.

Make sure you do not have any disks or diskettes in the computer system that are not write-protected—you may inadvertently write on a disk rather than read it during the early stages of testing.

Can a sector be written to the disk?

Using DDT, set up to write to a particular disk, track, and sector. Remove any write protection that you put on the target disk during earlier testing. You can either use the testbed's ripple pattern or fill the disk buffer area with a distinctive pattern. Write this data onto the disk, fill the buffer area with a *different* pattern, and read in the sector that you wrote. Check that the disk buffer gets changed back to the pattern written to the disk.

Does the driver display error messages correctly?

Rather than deliberately damaging a diskette to create errors, use DDT to temporarily sabotage the disk driver's logic. Make it return each of the possible error codes in turn, checking each time that the correct error message is displayed.

For each error condition in turn, check that the disk driver performs the correct recovery action, including interacting with the user and offering the choice of retrying, ignoring the error, or aborting the program.

Live Testing a New BIOS

Given that the drivers have passed all of the testing outlined above, you are ready to pull all of the BIOS pieces together and build a CP/M image.

For your initial testing, disable the real time clock, and use simple, polled I/O for the console driver if you can. It is important to get *something* up and running as soon as possible, and it is easier to do this without possible side effects from interrupts.

Prepare a complete listing of the BIOS and plan to spend at least an hour checking through it. Take a dry run through the console and disk driver—if there are any serious bugs left in these two drivers, CP/M may not start up. Remember that once the BIOS cold boot code has been executed and control is handed over to the CCP, the BDOS will be requested to log in the system disk, and this involves reading in the disk's directory.

Pay special attention to checking some of the major data structures. Make certain that everything is at a reasonable place in memory; for example, if the last address used by the BIOS is greater than 0FFFFH, you will need to move the entire CP/M image down in memory.

Then build a system disk, load it into the machine, and press the RESET button. You should see the bootstrap sign on, then the BIOS, and after a pause of about one second, the `A>` prompt (or `0A>` if you have included the special feature that patches the CCP).

If you see both sign-on messages but do not get an `A>` prompt, a likely cause of the problem is in the disk drivers. Alternatively, the directory area on the disk may be full of random data rather than `0E5H`'s.

If you cannot see what is wrong with the system, you might try faking the disk drivers to return a 128-byte block of `0E5H`'s for each read operation. The CCP should then sign on.

Once you do have the `A>` prompt, you can proceed with the system checkout. Start by checking that the warm boot logic works. Type a `CONTROL-C`. There should be a slight pause, and the `A>` prompt should be output again.

Next, check that you can read the disk directory by using the `DIR` command. If you have an empty directory, you should get a `NO FILE` response. If you get strange characters instead, you either forgot to initialize the directory area or the disk parameter block is directing CP/M to the wrong part of the disk for the file directory. If the system crashes, there is a problem with the disk driver.

Check that you can write on the disk by entering the command `SAVE I TEST`. Then use the `DIR` command to confirm that file `TEST` shows up in the file directory. If it does, use the `ERA` command `ERA TEST` and do another `DIR` command to confirm that `TEST` has indeed been erased.

If `TEST` either does not show up on the disk or cannot be erased, then you have a problem with the disk driver `WRITE` routine.

Put a standard CP/M release diskette into drive B and use the `DIR` command to check that you can access the drive and display a disk directory. If you do, then load the `DDT` utility and exit from it by using a `G0` (`G`, zero) command. This further tests if the disk drivers are functioning correctly.

To test the deblocking logic (if you are using disks that require deblocking), use the command:

```
PIP A:=B:*.*[V]
```

This copies all files from drive B to drive A using the verify option. It is a particularly good test of the system, and if you have any problems with the high-level disk drivers and deblocking code, you will get a Verify Error message from PIP. You can also get this message if you have hardware problems with the computer's memory, so run a memory test if you cannot find anything obviously wrong with the deblocking algorithm.

To completely test the deblocking code, you need to use PIP to copy a file of text larger than the amount of memory available. Thus, you may have to create a large text file using a text editor just to provide PIP with test data.

With the disk driver functioning correctly, rebuild the system with the real time clock enabled. Bring up the new system and check that the ASCII time of day is

being updated in the configuration block; use DDT to inspect this in memory. Set the clock to the current time, let it run for five minutes, and see if it is still accurate. You may have to adjust one of the initialization time constants for the device that is providing the periodic interrupts for the clock.

Rebuild the system yet again, this time with the real interrupt-driven console input and the real console output routines. Check that the system comes up properly and that the initial forced-input startup string appears on the console.

Check that when you type characters on the keyboard they are displayed as you type them. If not, there could be a problem with either the CONIN or CONOUT routines. Experimentally type in enough characters to fill the input buffer. If the terminal's bell starts to sound, the interrupt service routine is probably not the culprit. Check the CONOUT routine again.

Check that the function key decode logic is working correctly. With the A> prompt displayed, press a function key. The CONIN driver should inject the correct function key string and it should appear on the terminal. For example, with the BIOS in Figure 8-10, pressing PF1 on the VT-100 terminal should produce this on the display:

```
A>Function Key1
Function?
A>
```

The CCP does not recognize "Function" as a legitimate command name, nor is there such a COM file—hence the question mark.

Using DDT, write a small program that outputs ESCAPE, "t" to the console, and check that the ASCII time of day string appears on the console. This checks that the escape sequence has been recognized.